

# Simultaneous Blind Deconvolution and Blind Demixing via Convex Programming

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Nov.8th, 2016

# Acknowledgements

Research in collaboration with:

Prof. Thomas Strohmer (UC Davis)

This work is sponsored by NSF-DMS and DARPA.



- Setup: blind deconvolution and demixing
- Convex relaxation and main result
- Numerics and idea of proof

# What is blind deconvolution?

## What is blind deconvolution?

Suppose we observe a function  $\mathbf{y}$  which is the convolution of two unknown functions, the blurring function  $\mathbf{f}$  and the signal of interest  $\mathbf{g}$ , plus noise  $\mathbf{w}$ . How to reconstruct  $\mathbf{f}$  and  $\mathbf{g}$  from  $\mathbf{y}$ ?

$$\mathbf{y} = \mathbf{f} * \mathbf{g} + \mathbf{w}.$$

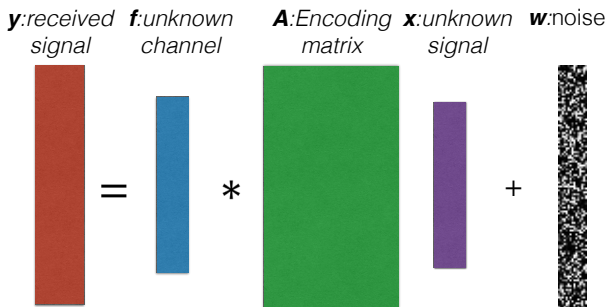
It is obviously a highly ill-posed **bilinear inverse** problem... but important in signal processing.

# Blind deconvolution in wireless communication

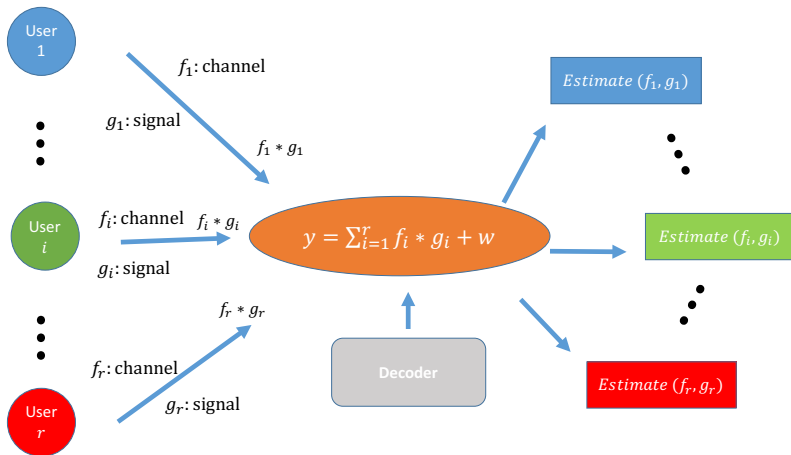
## Joint channel and signal estimation in wireless communication

Suppose that a signal  $\mathbf{x}$ , encoded by  $\mathbf{A}$ , is transmitted through an unknown channel  $\mathbf{f}$ . How to reconstruct  $\mathbf{f}$  and  $\mathbf{x}$  from  $\mathbf{y}$ ?

$$\mathbf{y} = \mathbf{f} * \mathbf{A}\mathbf{x} + \mathbf{w}.$$



# Blind deconvolution meets demixing?



# Blind deconvolution and blind demixing

We start from the original model

$$\mathbf{y} = \sum_{i=1}^r \mathbf{f}_i * \mathbf{g}_i + \mathbf{w}.$$

This is even more difficult than blind deconvolution, since this is a “mixture” of blind deconvolution problem

## More assumptions

- Each impulse responses  $\mathbf{f}_i$  has maximum delay spread  $K$ .

$$\mathbf{f}_i(n) = 0, \quad \text{for } n > K.$$

- $\mathbf{g}_i := \tilde{\mathbf{A}}_i \mathbf{x}_i$  is the signal  $\mathbf{x}_i \in \mathbb{C}^N$  encoded by matrix  $\tilde{\mathbf{A}}_i \in \mathbb{C}^{L \times N}$  with  $L > N$ .

# Model under subspace assumption

In the **frequency** domain,

$$\hat{\mathbf{y}} = \sum_{i=1}^r \hat{\mathbf{f}}_i \odot \hat{\mathbf{g}}_i + \mathbf{w} = \sum_{i=1}^r \text{diag}(\hat{\mathbf{f}}_i) \hat{\mathbf{g}}_i + \mathbf{w},$$

where “ $\odot$ ” denotes entry-wise multiplication. We assume  $\mathbf{y}$  and  $\hat{\mathbf{y}}$  are both of length  $L$ .

## Subspace assumption

Denote  $\mathbf{F}$  as the  $L \times L$  DFT matrix.

- Let  $\mathbf{h}_i \in \mathbb{C}^K$  be the first  $K$  nonzero entries of  $\mathbf{f}_i$  and  $\mathbf{B}_i$  be a low-frequency DFT matrix. There holds,

$$\hat{\mathbf{f}}_i = \mathbf{F} \mathbf{f}_i = \mathbf{B}_i \mathbf{h}_i.$$

- $\hat{\mathbf{g}}_i := \mathbf{A}_i \mathbf{x}_i$  where  $\mathbf{A}_i := \mathbf{F} \tilde{\mathbf{A}}_i$  and  $\mathbf{x}_i \in \mathbb{C}^N$ .



# Mathematical model

Finally, we end up with the following model,

## Model with subspace constraint

$$\mathbf{y} = \sum_{i=1}^r \text{diag}(\mathbf{B}_i \mathbf{h}_i) \mathbf{A}_i \mathbf{x}_i + \mathbf{w},$$

**Goal:** We want to recover  $(\mathbf{h}_i, \mathbf{x}_i)_{i=1}^r$  from  $(\mathbf{y}, \mathbf{B}_i, \mathbf{A}_i)_{i=1}^r$ .

**Remark:** The degree of freedom for unknowns:  $r(K + N)$ ; number of constraint:  $L$ . To make the solution identifiable, we require  $L \geq r(K + N)$  at least.

## Special case if $r = 1$

In particular, if  $r = 1$ , it is a blind deconvolution problem.

$$\mathbf{y} = \text{diag}(\mathbf{B}\mathbf{h})\mathbf{A}\mathbf{x} + \mathbf{w}.$$

# Nonconvex optimization?

## Naive approach?

We may want to try nonlinear least squares approach:

$$\min_{(\mathbf{u}_i, \mathbf{v}_i)} \left\| \sum_{i=1}^r \text{diag}(\mathbf{B}_i \mathbf{u}_i) \mathbf{A}_i \mathbf{v}_i - \mathbf{y} \right\|^2.$$

This gives a nonconvex objective function.

- May get stuck at local minima and no guarantees for recoverability.
- For  $r = 1$ , we have recovery guarantees by adding regularizers but not for  $r > 1$ .

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## Two-step convex approach

- (a) Lifting: convert nonconvex constraints to linear
- (b) Solving a SDP relaxation and hope to recover  $\{\mathbf{h}_i \mathbf{x}_i^*\}_{i=1}^r$

# Convex approach of demixing problem

## Step 1: lifting

Let  $\mathbf{a}_{i,l}$  be the  $l$ -th column of  $\mathbf{A}_i^*$  and  $\mathbf{b}_{i,l}$  be the  $l$ -th column of  $\mathbf{B}_i^*$ .

$$y_l = \sum_{i=1}^r (\mathbf{B}_i \mathbf{h}_i)_l \mathbf{x}_i^* \mathbf{a}_{i,l} + w_l = \sum_{i=1}^r \mathbf{b}_{i,l}^* \mathbf{h}_i \mathbf{x}_i^* \mathbf{a}_{i,l} + w_l,$$

Let  $\mathbf{X}_i := \mathbf{h}_i \mathbf{x}_i^*$  and define the linear operator  $\mathcal{A}_i : \mathbb{C}^{K \times N} \rightarrow \mathbb{C}^L$  as,

$$\mathcal{A}_i(\mathbf{Z}) := \{\mathbf{b}_{i,l}^* \mathbf{Z} \mathbf{a}_{i,l}\}_{l=1}^L = \{\langle \mathbf{Z}, \mathbf{b}_{i,l} \mathbf{a}_{i,l}^* \rangle\}_{l=1}^L.$$

Then, there holds

$$\mathbf{y} = \sum_{i=1}^r \mathcal{A}_i(\mathbf{X}_i) + \mathbf{w}.$$

- Advantage: linear constraints (convex constraints)
- Disadvantage: dimension increases

# Rank- $r$ matrix recovery

## Recast as rank- $r$ matrix recovery

We rewrite  $\mathbf{y} = \sum_{i=1}^r \text{diag}(\mathbf{B}_i \mathbf{h}_i) \mathbf{A}_i \mathbf{x}_i$  as

$$y_l = \left\langle \begin{bmatrix} \mathbf{h}_1 \mathbf{x}_1^* & \mathbf{0} & \cdots & \mathbf{0} \\ \mathbf{0} & \mathbf{h}_2 \mathbf{x}_2^* & \cdots & \mathbf{0} \\ \vdots & \vdots & \ddots & \vdots \\ \mathbf{0} & \mathbf{0} & \cdots & \mathbf{h}_r \mathbf{x}_r^* \end{bmatrix}, \begin{bmatrix} \mathbf{b}_{1,l} \mathbf{a}_{1,l}^* & \mathbf{0} & \cdots & \mathbf{0} \\ \mathbf{0} & \mathbf{b}_{2,l} \mathbf{a}_{2,l}^* & \cdots & \mathbf{0} \\ \vdots & \vdots & \ddots & \vdots \\ \mathbf{0} & \mathbf{0} & \cdots & \mathbf{b}_{r,l} \mathbf{a}_{r,l}^* \end{bmatrix} \right\rangle$$

- Find a **rank- $r$  block diagonal** matrix satisfying the linear constraints above.
- Finding such a rank- $r$  matrix is also an NP-hard problem.

## Nuclear norm minimization

Since this system is highly underdetermined, we hope to recover all  $\{\mathbf{Z}_i\}_{i=1}^r$  from

$$\min \sum_{i=1}^r \|\mathbf{Z}_i\|_* \quad \text{subject to} \quad \sum_{i=1}^r \mathcal{A}_i(\mathbf{Z}_i) = \mathbf{y}.$$

Once we obtain  $\{\hat{\mathbf{Z}}_i\}_{i=1}^r$ , we can easily extract the leading left and right singular vectors from  $\mathbf{Z}_i$  as the estimation of  $(\mathbf{h}_i, \mathbf{x}_i)$ .

**Key question:** does the solution to the SDP above really give  $\{\mathbf{h}_i \mathbf{x}_i^*\}_{i=1}^r$ ?  
What conditions are needed?

# State-of-the-art: blind deconvolution ( $r = 1$ )

## Nuclear norm minimization

Consider the convex envelop of  $\text{rank}(\mathbf{Z})$ : nuclear norm  $\|\mathbf{Z}\|_* = \sum \sigma_i(\mathbf{Z})$ .

$$\min \|\mathbf{Z}\|_* \quad \text{s.t.} \quad \mathcal{A}(\mathbf{Z}) = \mathcal{A}(\mathbf{X}).$$

Convex optimization can be solved within polynomial time.

## Theorem [Ahmed-Recht-Romberg 14]

Assume  $\mathbf{y} = \text{diag}(\mathbf{B}\mathbf{h})\mathbf{A}\mathbf{x}$ ,  $\mathbf{A} : L \times N$  is a Gaussian random matrix, and  $\mathbf{B} \in \mathbb{C}^{L \times K}$  is a partial DFT matrix,

$$\mathbf{B}^* \mathbf{B} = \mathbf{I}_K, \quad L \|\mathbf{B}\mathbf{h}\|_\infty^2 \leq \mu_h^2,$$

the above convex relaxation recovers  $\mathbf{X} = \mathbf{h}\mathbf{x}^*$  exactly with high probability if

$$C_0 \max(K, \mu_h^2 N) \leq \frac{L}{\log^3 L}.$$

## Theorem [Ling-Strohmer 15]

Each  $\mathbf{B}_i \in \mathbb{C}^{L \times K}$  partial DFT matrix with  $\mathbf{B}_i^* \mathbf{B}_i = \mathbf{I}_K$  and each  $\mathbf{A}_i$  is a Gaussian random matrix, i.e., each entry in  $\mathbf{A}_i \stackrel{\text{i.i.d.}}{\sim} \mathcal{N}(0, 1)$ . Let  $\mu_h^2$  be as defined in  $\mu_h^2 = L \max_{1 \leq i \leq r} \frac{\|\mathbf{B}_i \mathbf{h}_i\|_\infty^2}{\|\mathbf{h}_i\|^2}$ . If

$$L \geq C_{\alpha + \log r} r^2 \max\{K, \mu_h^2 N\} \log^3 L,$$

then the solution to convex relaxation satisfies

$$\hat{\mathbf{X}}_i = \mathbf{X}_i, \quad \text{for all } i = 1, \dots, r,$$

with probability at least  $1 - \mathcal{O}(L^{-\alpha+1})$ .



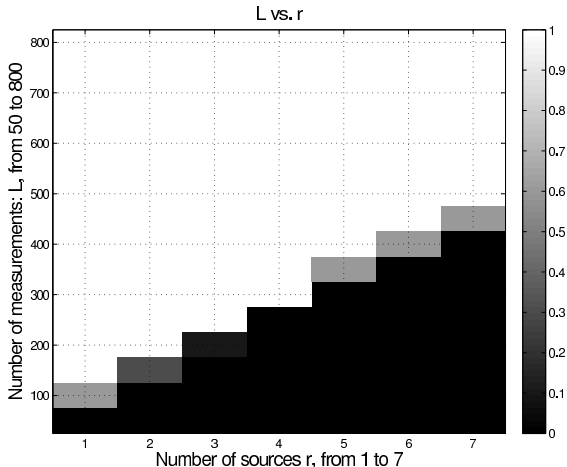
## Remark

- Our result is a generalization of Ahmed-Romberg-Recht's result to  $r > 1$ .
- $\mathbf{B}_j$  have other choices other than DFT matrix.
- Incoherence  $\mu_h^2$  does affect the result.

- Our result is a generalization of Ahmed-Romberg-Recht's result to  $r > 1$ .
- $\mathbf{B}_i$  have other choices other than DFT matrix.
- Incoherence  $\mu_h^2$  does affect the result.
- $r^2$  is not optimal. One group has claimed to reduce the sampling complexity from  $r^2$  to  $r$ .
- Empirically,  $\mathbf{A}_i$  can be other matrices besides Gaussian. However, no theories exist so far.

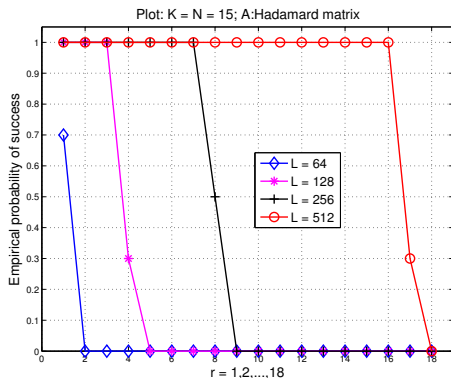
# Numerics: does $L$ really scales linearly with $r$ ?

$\mathbf{A}_i$  is chosen as Gaussian matrix. Here  $K = 30$  and  $N = 25$  are fixed.  
Black: failure; White: success.



# Numerics: does $L$ really scales linearly with $r$ ?

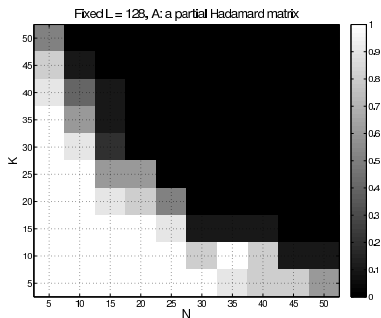
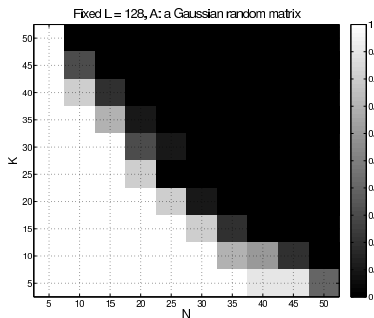
Choose  $\mathbf{A}_i = \mathbf{D}_i \mathbf{H}$  where  $\mathbf{H}$  is a partial Hadamard matrix ( $\pm 1$  orthogonal matrix).  $\mathbf{D}_i$  is a random  $\pm 1$  diagonal matrix.



Number of unknowns  $r(K + N) = 16 \cdot 30 = 480$  is slightly smaller than the number of constraints 512.

# Numerics: $L$ scales linearly with $K + N$ .

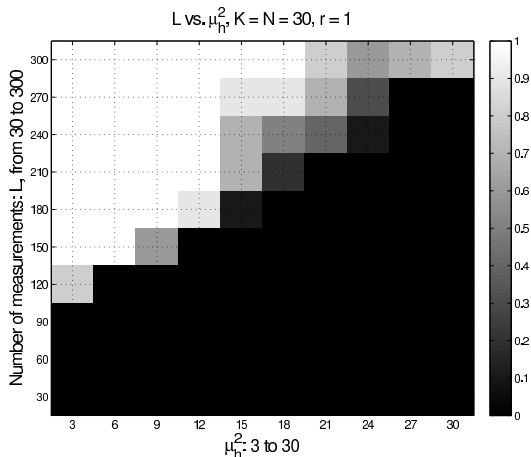
Numerical simulations verify our theory that  $L \approx \mathcal{O}(r(K + N))$  gives exact recovery. Here  $r = 2$  and  $L \approx 1.5r(K + N)$ .



# Numerics: $L$ vs. $\mu_h^2$

We observe strong linear correlation between minimal required  $L$  and  $\mu_h^2$ :

$$L \propto \max_{1 \leq i \leq r} \frac{\|B_i h_i\|_\infty^2}{\|h_i\|^2}.$$



# Stability theorem

In reality measurements are noisy. Hence, suppose that  $\hat{\mathbf{y}} = \mathbf{y} + \mathbf{w}$  where  $\mathbf{w}$  is noise with  $\|\mathbf{w}\| \leq \eta$ .

$$\min \sum_{i=1}^r \|\mathbf{z}_i\|_* \quad \text{subject to} \quad \left\| \sum_{i=1}^r \mathcal{A}_i(\mathbf{z}_i) - \hat{\mathbf{y}} \right\| \leq \eta. \quad (1)$$

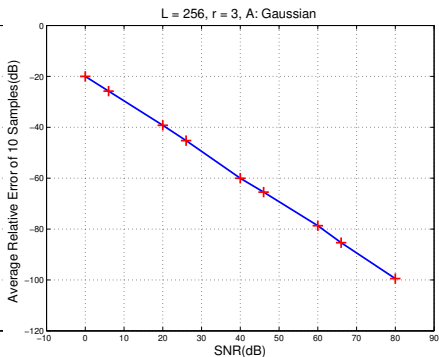
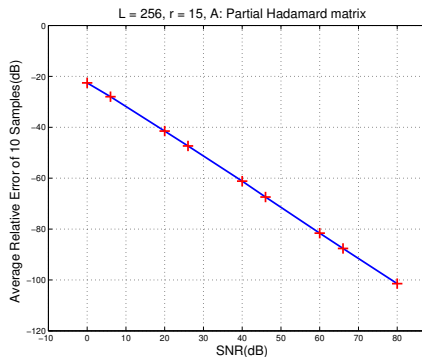
## Theorem

Assume we observe  $\hat{\mathbf{y}} = \mathbf{y} + \mathbf{w} = \sum_{i=1}^r \mathcal{A}_i(\mathbf{x}_i) + \mathbf{w}$  with  $\|\mathbf{w}\| \leq \eta$ . Then, the minimizer  $\{\hat{\mathbf{x}}_i\}_{i=1}^r$  satisfies

$$\sqrt{\sum_{i=1}^r \|\hat{\mathbf{x}}_i - \mathbf{x}_i\|_F^2} \leq Cr \sqrt{\max\{K, N\}} \eta.$$

with probability at least  $1 - \mathcal{O}(L^{-\alpha+1})$ .

We see that the relative error is **linearly** correlated with the noise in dB. Approximately, 10 units of increase in SNR leads to the *same* amount of decrease in relative error (in dB).





# Sketch of proof

Let's consider the noiseless version,

$$\min \sum_{i=1}^r \|\mathbf{z}_i\|_*, \quad \text{subject to } \sum_{i=1}^r \mathcal{A}_i(\mathbf{z}_i) = \mathbf{y}.$$

Difficulties:

- $\mathbf{X}_i$  is asymmetric.
- How to deal with block diagonal structure?

## Two-step proof

- Find a sufficient condition for exact recovery
- Construct an approximate dual certificate via golfing scheme

## Three key ingredients to achieve exact recovery

- 1 Local isometry property on  $T_i$

$$\max_{1 \leq i \leq r} \|\mathcal{P}_{T_i} \mathcal{A}_i^* \mathcal{A}_i \mathcal{P}_{T_i} - \mathcal{P}_{T_i}\| \leq \frac{1}{4}$$

where  $T_i = \{\mathbf{h}_i \mathbf{h}_i^* \mathbf{Z} + (\mathbf{I} - \mathbf{h}_i \mathbf{h}_i^*) \mathbf{Z} \mathbf{x}_i \mathbf{x}_i^*\}$ .

- 2 Local incoherence property

$$\max_{i \neq j} \|\mathcal{P}_{T_i} \mathcal{A}_i^* \mathcal{A}_j \mathcal{P}_{T_j}\| \leq \frac{1}{4r}$$

- 3 Existence of an approximate dual certificate, which is achieved via the celebrated golfing scheme). Find a  $\boldsymbol{\lambda} \in \mathbb{C}^L$  such that for all  $1 \leq i \leq r$ ,

$$\|\mathcal{P}_{T_i}(\mathcal{A}_i^* \boldsymbol{\lambda}) - \mathbf{h}_i \mathbf{x}_i^*\|_F \leq \frac{1}{5r\gamma}, \quad \|\mathcal{P}_{T_i^\perp}(\mathcal{A}_i^* \boldsymbol{\lambda})\| \leq \frac{1}{2}$$

where  $\gamma := \max\{\|\mathcal{A}_j\|\}$ .

# Conclusion and future work

- Can we derive a theoretical bound that scales linearly in  $r$ , rather than quadratic in  $r$  as our current theory? (It may have been solved!)
- Is it possible to develop satisfactory theoretical bounds for deterministic matrices  $\mathbf{A}_i$ ?
- Fast algorithms: extend our nonconvex optimization framework to this blind-deconvolution-blind-demixing scenario.
- Can we develop a theoretical framework where the signals  $\mathbf{x}_i$  belong to some non-linear subspace, e.g. for sparse  $\mathbf{x}_i$ ?
- **See details in our paper:** *Blind Deconvolution Meets Blind Demixing: Algorithms and Performance Bounds.* [arXiv:1512.07730](https://arxiv.org/abs/1512.07730).