Homework is due by **11pm of Oct 14**. Send by email to both "regev" (under the cs.nyu.edu domain) and "des480" (under the nyu.edu domain) with subject line "CSCI-GA 3210 Homework 5" and name the attachment "YOUR NAME HERE HW5.tex/pdf". There is no need to print it. Start early!

- (4 points) (More indistinguishability) For a probability distribution D over Ω and positive integer m, let D^m denote the product distribution over Ω^m, obtained by drawing a tuple of m independent samples from D. Let X = {X_n} and Y = {Y_n} be ensembles of distributions that are efficiently sampleable (in PPT), and let m(n) = poly(n). Prove that if X ≈ Y, then {X_n^{m(n)}} ≈ {Y_n^{m(n)}}. (Where do you use that X_n, Y_n are efficiently sampleable?)
- 2. (5 points) (*Prediction vs distinguishing*) A function $h : \{0,1\}^* \to \{0,1\}$ is *hard-core* for a function f if for all PPT algorithms A,

$$\Pr_{x \leftarrow \{0,1\}^n} [\mathcal{A}(f(x)) = h(x)] \le \frac{1}{2} + \operatorname{negl}(n) .$$

Show that this definition is equivalent to requiring that

$$(f(U_n), h(U_n)) \stackrel{c}{\approx} (f(U_n), U_1),$$

where U_n is a uniform *n*-bit string, and U_1 is a uniform bit. Simplify the right hand side when *f* is a *permutation* (i.e., a bijection). Once you're done, I recommend reading Goldreich's Section 3.3.5

- 3. (*Hard core.*)¹ Prove or disprove (giving the simplest counterexample you can find) the following statements. In constructing a counterexample, you may assume the existence of another OWF / PRG.
 - (a) (1 point) If an efficiently-computable function f has a hard-core predicate h, then f is one-way.
 - (b) (3 points) If an efficiently-computable injective (one-to-one) function f has a hard-core predicate h, then f is one-way.
- 4. (2 points) (Pseudorandom functions⁴) We would like to extend the definition of a pseudorandom generator so that its output length is exponential. Can you think of a definition that makes sense?

¹A question from Peikert's class