

# Iterative Solvers

## Last Time Eigenvalue calculations

① Power method:

$$A^k w \approx \lambda_1^k v_1, \quad |\lambda_1| > |\lambda_2| > \dots$$

converges at a rate  $\sim O\left(\left|\frac{\lambda_2}{\lambda_1}\right|\right)$

② To increase convergence rate: shift

$A - sI$  has eigenvalues equal to  $\lambda_1 - s, \lambda_2 - s, \dots$

pick  $s$  to increase convergence rate.

Ex:  $\lambda_j = 1, 2, \dots, 8, 9, 10$

$\lambda_1 = 10, \lambda_2 = 9$

Power method converges at a rate of  $\left(\frac{9}{10}\right)^k$ .

Shift by  $s=1, \lambda_1=9, \lambda_2=8, \dots, \lambda_{10}=0$

$$\text{rate} = \left(\frac{8}{9}\right)^k < \left(\frac{9}{10}\right)^k$$

Best shift: shift such that  $|\lambda_1 - s| = |\lambda_{10} - s|$

$s = \frac{10+0}{2} = 5 \Rightarrow$

$\lambda_1 = 5$

$\lambda_2 = 4$

$\vdots$

$\lambda_{10} = -4$

$$\text{rate} = \left(\frac{4}{5}\right)^k = .8^k$$

(Eigenvectors did not change)

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③ To find interior eigenvalues, the inverse iteration with shift is required:

$$(A-sI)^{-1} \text{ has eigenvalues } \frac{1}{\lambda_1-s}, \dots, \frac{1}{\lambda_{10}-s} \dots$$

Ex: for  $s = 3.2$ ,  $\tilde{\lambda}_j$  are

$$\tilde{\lambda}_1 = \frac{1}{10-3.2}, \tilde{\lambda}_2 = \frac{1}{9-3.2} \dots \tilde{\lambda}_{10} = \frac{1}{1-3.2}$$

The largest absolute value  $\tilde{\lambda}$  is  $\tilde{\lambda}_8 = \frac{1}{3-3.2} = \frac{1}{-.8} = -5$ .

$$\tilde{\lambda}_7 = \frac{1}{.8} = 1.25$$

$$\tilde{\lambda}_9 = \frac{1}{-1.2} = -\frac{5}{6}$$

convergence rate:  $O\left(\left(\frac{1.25}{5}\right)^k\right)$ .

$\Rightarrow \tilde{\lambda}_j$ 's can be estimated with power method + deflation ...

How do we ~~iterate~~ compute  $(A-sI)^{-1} \vec{w}^{(k)}$ ?

$\Rightarrow$  solve  $(A-sI) \vec{w}^{(k+1)} = \vec{w}^{(k)}$

Intro: Iterative solver.

cost of Gaussian elimination:  $O(n^3)$

" " matrix-vector mult:  $O(n^2)$   $\leftarrow$  maybe

we can iterate

$k$  times

$\Rightarrow O(kn^2)$ .

Another reason we want iterative solvers: storage.

$$A = \begin{pmatrix} 2 & -1 & & & \\ -1 & 2 & \dots & & \\ & \dots & \dots & \dots & \\ & & & -1 & 2 & -1 \\ & & & & -1 & 2 \end{pmatrix}$$

Storage cost for A:  $O(3n)$

$A^{-1}$  is dense! (check in Matlab for yourself)

storage:  $O(n^2)$ . ~~this~~

If A is  $n \times n$ , and contains  $O(n)$  or  $O(n \log n)$  entries, we say that it is sparse.

### Simple Iteration

Suppose that you have a matrix M such that  $M^{-1}A \approx I$  and  $M\vec{x} = \vec{y}$  is easy to solve.

Then: To solve  $A\vec{x} = \vec{b}$ :

~~start~~

Start with guess  $\vec{x}_0$

Compute residual  $\vec{r}_0 = \vec{b} - A\vec{x}_0$

Solve  $M\vec{z}_0 = \vec{r}_0$ , then  $\vec{z}_0 = M^{-1}(\vec{b} - A\vec{x}_0)$

$$\approx M^{-1}\vec{b} - \vec{x}_0$$
~~$$\approx M^{-1}(\vec{b} - A\vec{x}_0)$$~~

$$\approx \vec{e}_0 = A^{-1}\vec{b} - \vec{x}_0$$

set  $\vec{x}_1 = \vec{x}_0 + \vec{z}_0$

And repeat:

$k=1, 2, \dots$

$$\text{Set } \vec{x}_k = \vec{x}_{k-1} + \vec{z}_{k-1}$$

$$\text{Compute } \vec{r}_k = \vec{b} - A\vec{x}_k$$

If  $\|\vec{r}_k\|$  is small, stop

$$\text{Else Solve } M\vec{z}_k = \vec{r}_k$$

$M$  is referred to as a left preconditioner.

How do we choose  $M$ ?

Ex:  $M =$  Lower triangle of  $A$ . (Gauss-Seidel iterations)

Easy to solve,  $O(n^2)$  flops by Fwd substitute

Convergence of Iterative Methods

Consider the matrix splitting:  $A = M - N$ .

$$A\vec{x} = \vec{b}$$

$$(M - N)\vec{x} = \vec{b}$$

$$\vec{x} = M^{-1}N\vec{x} + M^{-1}\vec{b}$$

← similar to fixed point iteration with  $\vec{q}(\vec{x}) = M^{-1}N\vec{x} + M^{-1}\vec{b}$

This is, in fact, equivalent to the earlier iteration

method:

$$\vec{x}_k = \vec{x}_{k-1} + \vec{z}_k \quad \left( \begin{aligned} M \vec{z}_k &= \vec{r}_k \\ &= \vec{b} - A \vec{x}_k \end{aligned} \right)$$

$$= \vec{x}_{k-1} + M^{-1}(\vec{b} - A \vec{x}_{k-1})$$

$$= (\mathbf{I} - M^{-1}A) \vec{x}_{k-1} + M^{-1}\vec{b}$$

and since

$$A = M - N$$

$$\underbrace{M^{-1}N}_{\text{same as before}} = \mathbf{I} - M^{-1}A = M^{-1}(\underbrace{M - A}_N)$$

So when does  $\vec{x}_k = (\mathbf{I} - M^{-1}A) \vec{x}_{k-1}$  converge?

At step  $k$ , the error is  $\vec{e}_k = A^{-1}\vec{b} - \vec{x}_k$

$$\vec{x}_k = \vec{x}_{k-1} + \vec{z}_{k-1}$$

$$\underbrace{-\vec{x}_k + A^{-1}\vec{b}}_{\vec{e}_k} = \underbrace{-\vec{x}_{k-1} + A^{-1}\vec{b}}_{\vec{e}_{k-1}} + \underbrace{\vec{z}_{k-1}}_{M^{-1}\vec{r}_{k-1}}$$

$$\vec{e}_k = \vec{e}_{k-1} + M^{-1}\vec{r}_{k-1}$$

$$= \vec{e}_{k-1} + M^{-1}A \vec{e}_{k-1}$$

$$= (\mathbf{I} - M^{-1}A) \vec{e}_{k-1}$$

$$\begin{aligned} \vec{r}_k &= \vec{b} - A \vec{x}_k \\ &= A(A^{-1}\vec{b} - \vec{x}_k) \\ &= A \vec{e}_k \end{aligned}$$

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This implies that

$$\vec{e}_k = (\mathbf{I} - \mathbf{M}^{-1}\mathbf{A})^k \vec{e}_0$$

$$\Rightarrow \|\vec{e}_k\| \leq \|(\mathbf{I} - \mathbf{M}^{-1}\mathbf{A})^k\| \|\vec{e}_0\|$$

any vector norm and  
the induced matrix norm.

Theorem:  $\|\vec{e}_k\| \rightarrow 0$  and  $\vec{x}_k \rightarrow \vec{x}$  if and only if

$$\lim_{k \rightarrow \infty} \|(\mathbf{I} - \mathbf{M}^{-1}\mathbf{A})^k\| = 0$$

for every initial  $\vec{x}_0$

(Unsurprising)

Example Recall  $\|A\|_2$ .

$$\|A\|_2 = \sqrt{\max_{\lambda} (A^T A)}$$

$$\Rightarrow \|A\|_2 = \max |\lambda_j|$$

Therefore we need the eigenvalues of  $\mathbf{I} - \mathbf{M}^{-1}\mathbf{A}$  to be  $\leq 1$  in absolute value.

Definition Spectral radius:

For an  $n \times n$  matrix  $G$ ,

$$\rho(G) = \max_j |\lambda_j|$$

Theorem: Let  $G$  be  $n \times n$  matrix. Then  $\lim_{k \rightarrow \infty} G^k = 0$   
if and only if  $\rho(G) < 1$ .

Idea: If  $G = VDV^{-1}$ , then  $G^k = VD^kV^{-1}$   
 $\downarrow$   
0 iff elements  $< 1$ .

$\Rightarrow$  for convergence, we need  $\rho(I - M^{-1}A) < 1$ .

(and if  $\|(I - M^{-1}A)^k\| \rightarrow 0$ , then  $\rho(I - M^{-1}A) < 1$ ).

Special case: symmetric positive definite matrix

No preconditioner, consider only multiply of  $A$ , and the

simple iteration:

$$\begin{aligned} \vec{x}_0 &= \vec{x}_0 \\ \vec{r}_0 &= b - A\vec{x}_0 = A\vec{e}_0 \\ \vec{x}_1 &= \vec{x}_0 + C_0 \vec{r}_0 \\ &= \vec{x}_0 + C_0 (b - A\vec{x}_0), \quad \vec{r}_1 = b - A\vec{x}_1 \\ \vec{x}_2 &= \vec{x}_1 + C_1 \vec{r}_1 \\ &= \vec{x}_0 + C_0 (b - A\vec{x}_0) + C_1 (b - A\vec{x}_1) \\ &= \vec{x}_0 + C_0 (b - A\vec{x}_0) + C_1 (b - A(\vec{x}_0 + C_0 (b - A\vec{x}_0))) \end{aligned}$$

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# Conjugate Gradient for RSPD matrix

Def:

$\vec{u}, \vec{v}$  are conjugate w.r.t. to  $A$  if

$$(\vec{u}, A\vec{v}) = 0 \quad (\text{the } A\text{-inner product})$$

Note  $(\vec{u}, A\vec{v})$  is only an inner product if  $A$  is symmetric and positive definite.

Let  $\{\vec{p}_1, \dots, \vec{p}_n\}$  be a set of  $n$  mutually  $A$ -conjugate vectors. Then they form a basis for  $\mathbb{R}^n$  (if  $A$  is invertible).

Any  $\vec{x}$  can be written as

$$\vec{x} = \sum \alpha_j \vec{p}_j.$$

to solve  $A\vec{x} = \vec{b}$ :

Therefore:

$$A\vec{x} = \sum \alpha_j A\vec{p}_j = \sum \alpha_j A\vec{p}_j$$

$$(\vec{p}_k, A\vec{x}) = \sum \alpha_j (p_k, A p_j) = \sum \alpha_j (p_k, A p_j) (p_k, \vec{b})$$

$$= \alpha_k (p_k, A p_k) = (p_k, \vec{b})$$

$$\Rightarrow \alpha_k = \frac{(p_k, \vec{b})}{\|p_k\|_A}$$

Direct solution.



To do this iteratively, is there a way to select  $\vec{p}_1, \dots, \vec{p}_n$  so that only a few inner products need to be calculated? (as an approximation to  $\vec{x}$  s.t.  $\vec{r} = \vec{b} - A\vec{x}$  is small).

Observation: The solution to  $A\vec{x} = \vec{b}$  is also

the unique minimizer of the function:

$$f(\vec{x}) = \frac{1}{2}(\vec{x}, A\vec{x}) - (\vec{x}, \vec{b})$$

(I.e.  $\frac{df}{dx_j} = 0$  when  $A\vec{x} = \vec{b}$ ).

Newton would suggest moving in the direction of  $-\nabla f(\vec{x}_0)$ .  
minimize

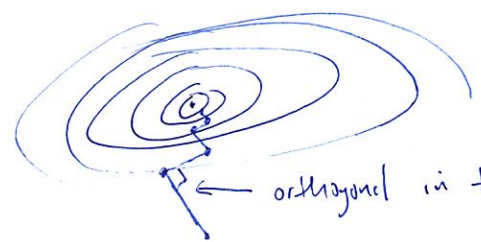
$$\nabla f(\vec{x}) = A\vec{x} - \vec{b}$$

$$\text{Let } \vec{p}_0 = \vec{b} - A\vec{x} = \vec{r}_0 = -\nabla f(\vec{x}_0).$$

We want the following vectors  $\vec{p}_1, \vec{p}_2, \dots$  to be A-conjugate:

$$\text{Let } \vec{r}_k = \vec{b} - A\vec{x}_k = -\nabla f(\vec{x}_k)$$

$$\Rightarrow \vec{p}_k = \vec{r}_k - \sum_{l < k} \frac{(\vec{p}_l, A\vec{r}_k)}{(\vec{p}_l, A\vec{p}_l)} \vec{p}_l$$



orthogonal in the A-inner product.

“A-inner product”  
 Gram-Schmidt

In the direction of  $\vec{p}_k$ , the next iterate is:

$$\vec{x}_{k+1} = \vec{x}_k + \alpha_k \vec{p}_k$$

$$\text{with } \alpha_k = \frac{(\vec{p}_k, \vec{b})}{(\vec{p}_k, A\vec{p}_k)} = \frac{(\vec{p}_k, \vec{r}_{k-1} + A\vec{x}_{k-1})}{(\vec{p}_k, A\vec{p}_k)} = \frac{(\vec{p}_k, \vec{r}_{k-1})}{(\vec{p}_k, A\vec{p}_k)}$$

since  $(\vec{p}_k, A\vec{x}_{k-1})$  are A-conjugate.

### Convergence Rate

Thm Let  $\vec{e}_k$  be the error on step  $k$  of CG. Then

$$\frac{\|\vec{e}_k\|_A}{\|\vec{e}_0\|_A} \leq 2 \left( \frac{\sqrt{K}-1}{\sqrt{K}+1} \right)^k$$

$$K = \frac{\lambda_{\max}}{\lambda_{\min}} = \text{condition number}$$

(for SPD matrices)

Conjugate Gradient is a Krylov Method

It minimizes some norm of  $\|\vec{r}_k\| \leftarrow$  in fact the A-norm.  
using  $\{\vec{r}_0, A\vec{r}_0, A^2\vec{r}_0, \dots\}$

For nonsymmetric, arbitrary systems  $A\vec{x} = \vec{b}$ , the GMRES algorithm minimizes the 2 norm  $\|\vec{r}_k\|$ .

Idea: solve  $A\vec{x} = \vec{b}$  by choosing  
 an approximate iterate  $\vec{x}_k$  that is of  
 the form  $\vec{x}_0 + \{\vec{r}_0, A\vec{r}_0, \dots, A^{k-1}\vec{r}_0\}$  to

minimize  $\vec{r}_k = \vec{r}_0 - \sum_1^k c_j A^j \vec{r}_0$  in the 2-norm.

(Use Gram-Schmidt) and an orthogonal basis  
 for  $\{\vec{r}_0, A\vec{r}_0, A^2\vec{r}_0, \dots\}$