Mehryar Mohri Advanced Machine Learning 2021 Courant Institute of Mathematical Sciences Homework assignment 2 April 20, 2021 Due: May 04, 2021

## A. Structural Risk Minimization

As discussed in class, the Structural Risk Minimization (SRM) technique is based on a hypothesis set  $\mathcal{H}$  defined as a countable union of hypothesis sets  $\mathcal{H}_n$  with finite VC-dimension or favorable Rademacher complexity. In this problem, we study several questions related to such countable union hypothesis sets.

- 1. Let  $\mathcal{H} = \bigcup_{n=1}^{+\infty} \{h_n\}$  be a countable hypothesis set and assume that the target labeling function is in  $\mathcal{H}$ . In the standard statistical learning scenario, the learner receives an i.i.d. sample that he uses to train an algorithm and return a predictor. Here, suppose instead that the learner can request more labeled samples drawn i.i.d., as needed. Consider the following algorithm: starting from t = 1, at each round t, sample  $m_t = \frac{1}{\epsilon} \log \frac{1}{\delta_t}$  labeled points; if  $h_t$  is consistent with  $m_t$ , return  $h_t$  and stop.
  - (a) Prove that the algorithm terminates.

Solution: Since the Bayes classifier  $f^*$  is in  $\mathcal{H}$ , there exists t such that  $f^* = h_t$ , thus the algorithm terminates at most after t rounds.

(b) Fix  $\epsilon, \delta > 0$  and choose  $\delta_t = \frac{\delta}{2t^2}$ . Show that with probability  $1 - \delta$ , the algorithm returns a hypothesis with error at most  $\epsilon$ . Suppose we use the samples obtained from previous rounds to test consistency, then, what is the maximum number of samples needed by the algorithm?

Solution: The probability that the algorithm stops at round twhile  $h_t$  has error  $\epsilon$  is  $\mathbb{P}[h_t \text{ consistent}|R(h_t) \ge \epsilon] \le (1-\epsilon)^{m_t} \le e^{-\epsilon m_t} = \delta_t$ . Thus, by the union bound,

$$\mathbb{P}[\exists t \ge 1 \colon h_t \text{ consistent} | R(h_t) \ge \epsilon] \le \sum_{t=1}^{+\infty} \delta_t = \frac{\delta}{2} \sum_{t=1}^{+\infty} \frac{1}{t^2} = \frac{\delta}{2} \frac{\pi^2}{6} \le \delta_t$$

Let  $t^*$  be the time at which the algorithm terminates.  $t^*$  is upper bounded by the index t such that  $h_t = f^*$ . If we reuse samples, at most  $\frac{1}{\epsilon} \log \frac{2t^{*2}}{\delta}$  points are needed overall.

(c) Can you generalize these results to the case where  $\mathcal{H} = \bigcup_{n=1}^{+\infty} \mathcal{H}_n$ with  $\operatorname{VCdim}(\mathcal{H}_n) = d_n < +\infty$ ?

Solution: Same algorithm, except at round t a consistent hypothesis in  $\mathcal{H}_t$  is sought. Assume that the ordering of  $\mathcal{H}_n$  is such that  $\mathcal{H}_n \subset \mathcal{H}_{n+1}$ . At each round t, select a sample  $S_{m_t}$  of size  $m_t$  and return  $h_t \in \mathcal{H}_t$  if it is consistent with  $S_{m_t}$ . To derive the error bound, let  $\delta_t = \frac{\delta}{2t^2}$  and let  $m_t = O\left(\frac{d_t}{\epsilon} \log \frac{1}{\delta_{t\epsilon}}\right)$  and observe that:

$$\mathbb{P}\left(R_{\mathcal{D}}(h_t) > \epsilon\right) \leq \mathbb{P}\left(\bigcup_{t=0}^{\infty} \left\{ \exists h \in \mathcal{H}_t : \widehat{R}_{S_{m_t}}(h) = 0, R_{\mathcal{D}}(h) > 0 \right\} \right)$$
$$\leq \sum_{t=1}^{\infty} \delta_t$$
$$= \frac{\delta}{2} \sum_{t=1}^{\infty} \frac{1}{t^2}$$
$$\leq \delta.$$

- 2. Suppose S is an infinite set that can be fully shattered by  $\mathcal{H}$ . We wish to show that  $\mathcal{H}$  cannot be written as a countable union  $\mathcal{H} = \bigcup_{n=1}^{+\infty} \mathcal{H}_n$  with  $\operatorname{VCdim}(\mathcal{H}_n) = d_n < +\infty$ .
  - (a) Show that we can define a family of subsets  $(X_n)_{n\geq 1}$  such that  $|X_n| = d_n + 1$  and  $X_n \subseteq S \bigcup_{1 \le k \le n-1} X_k$ .

Solution: This is straightforward since S is an infinite sample and since  $d_n$  is finite for any  $n \ge 1$ .

(b) Show that for any  $n \ge 1$ , there exists a labeling  $X_n^l$  that cannot be obtained using  $\mathcal{H}_n$ .

Solution: This follows directly the definition of the VC-dimension: no set of size  $d_n + 1$  can be fully shattered by  $\mathcal{H}_n$ .

(c) Consider the labeling  $X^l$  of  $X = \bigcup_{n=1}^{+\infty} X_n$  obtained using all the  $X_n^l$ s. Show that no labeling of S using  $\mathcal{H}$  can be consistent with  $X^l$ . Conclude that that  $\mathcal{H}$  cannot be written as a countable union  $\mathcal{H} = \bigcup_{n=1}^{+\infty} \mathcal{H}_n$  with  $\operatorname{VCdim}(\mathcal{H}_n) = d_n < +\infty$ .

Solution: Note that, by definition, all  $X_n$ s are disjoint. Thus, the labeling  $X^l$  obtained from all  $X_n^l$ s is well defined. Let Y be a labeling of T consistent with  $X^l$ . Then, for any  $n \ge 1$ ,  $Y_{|X_n}$  is a labeling of  $X_n$  matching  $X_n^l$  and thus Y is not in  $\mathcal{H}_n$ . Since Y is not in  $\mathcal{H}_n$  for any  $n \ge 1$ , it is not in  $\mathcal{H}$ . This shows that the assumption that  $\mathcal{H}$  cannot be written as a countable union  $\mathcal{H} = \bigcup_{n=1}^{+\infty} \mathcal{H}_n$  with VCdim $(\mathcal{H}_n) = d_n < +\infty$  does not hold.

3. Suppose you only know an upper bound  $\alpha_n$  on VCdim $(\mathcal{H}_n) = d_n < +\infty$  with  $\sum_{n=1}^{+\infty} e^{-\alpha_n} < +\infty$ . Give a generalization bound for the SRM-type algorithm defined by

$$f^* = \operatorname*{argmin}_{k \ge 1, h \in \mathcal{H}_k} \widehat{R}_S + \sqrt{\frac{32\alpha_k \log(em)}{m}},$$

for a sample S of size m.

Solution: Let 
$$F_k(h) = \widehat{R}_S + \sqrt{\frac{32\alpha_k \log(em)}{m}}$$
. Then using  $\mathcal{H} = \bigcup_{k=1}^{+\infty} \mathcal{H}_k$ 
$$\mathbb{P}\left(\sup_{h \in \mathcal{H}} R(h) - F_{k(h)}(h) - \sqrt{\frac{2dk(h)\log em/d_{k(h)}}{m}} > \epsilon\right)$$

can be bounded as follows:

$$\leq \sum_{k=1}^{\infty} \mathbb{P}\left(\sup_{h \in \mathcal{H}_{k}} R(h) - F_{k}(h) - \sqrt{\frac{2dk \log em/d_{k}}{m}} > \epsilon\right)$$

$$= \sum_{k=1}^{\infty} \mathbb{P}\left(\sup_{h \in \mathcal{H}_{k}} R(h) - F_{k}(h) - \widehat{R}_{S}(h) - \sqrt{\frac{2dk \log em/d_{k}}{m}} > \epsilon + \sqrt{\frac{32\alpha_{k} \log(em)}{m}}\right)$$

$$\leq \sum_{k=1}^{\infty} \exp\left(-2m\left(\epsilon + \sqrt{\frac{32\alpha_{k} \log(em)}{m}}\right)^{2}\right)$$

$$\leq \sum_{k=1}^{\infty} \exp\left(-2m\epsilon^{2}\right) \exp\left(-a_{k} \log m\right)$$

$$\leq Ce^{-2m\epsilon^{2}}.$$

Applying similar steps and recalling that  $f^*$  is the minimizer of  $\widehat{R}_S$  +

$$\sqrt{\frac{32\alpha_k \log(em)}{m}}, \text{ we can show that}$$
$$\mathbb{P}\left(\sup_{h \in \mathcal{H}} F_{k(f^\star)}(f^\star) - R(h^\star) - \sqrt{\frac{32\alpha_k(h^\star)\log(em)}{m}} - \sqrt{\frac{2dk(h^\star)\log em/d_{k(h^\star)}}{m}} > \frac{\epsilon}{2}\right)$$
$$\leq e^{\frac{-m\epsilon^2}{2}}.$$

Combining the results above and the union bound provides the generalization bound with  $\delta = (1+C)e^{\frac{-m\epsilon^2}{2}}$ .

## **B.** Learning kernels

Let  $\mathcal{K}$  be the family of all Gaussian kernels defined over  $\mathbb{R}^N$ :

$$\mathcal{K} = \left\{ K_{\gamma} \colon K_{\gamma}(\mathbf{x}, \mathbf{x}') = e^{-\gamma \|\mathbf{x} - \mathbf{x}'\|^2}, \forall \mathbf{x}, \mathbf{x}' \in \mathbb{R}^N, \gamma > 0 \right\}.$$

Consider the hypothesis set defined via the reproducing kernel Hilbert space of the kernels in  $\mathcal{K}$ :

$$\mathcal{H} = \Big\{ h \colon h \in \mathbb{H}_K, K \in \mathcal{K}, \|h\|_{\mathbb{H}_K} \le 1 \Big\}.$$

1. Let  $S = (\mathbf{x}_1, \dots, \mathbf{x}_m)$  be a sample of size m. Show that  $\widehat{\mathfrak{R}}_S(\mathfrak{H}) = \frac{1}{m} \mathbb{E}_{\boldsymbol{\sigma}} \left[ \sqrt{\sup_{\gamma>0} \boldsymbol{\sigma}^\top \mathbf{K}_{\gamma} \boldsymbol{\sigma}} \right]$ , where  $\mathbf{K}_{\gamma}$  is the Gram matrix of kernel  $K_{\gamma}$  for the sample S.

Solution:

$$\begin{aligned} \widehat{\mathfrak{R}}_{S}(\mathcal{H}) &= \frac{1}{m} \operatorname{\mathbb{E}} \left[ \sup_{\substack{h \in \mathbb{H}_{K}, \|h\|_{\mathbb{H}_{k}} \leq 1 \\ K \in \mathcal{K}}} \sum_{i=1}^{m} \sigma_{i} \langle h, \Phi_{K}(x_{i}) \rangle \right] \\ &= \frac{1}{m} \operatorname{\mathbb{E}} \left[ \sup_{\substack{h \in \mathbb{H}_{K}, \|h\|_{\mathbb{H}_{k}} \leq 1 \\ K \in \mathcal{K}}} \left| \langle h, \sum_{i=1}^{m} \sigma_{i} \Phi_{K}(x_{i}) \right| \right] \\ &= \frac{1}{m} \operatorname{\mathbb{E}} \left[ \sup_{K \in \mathcal{K}} \left\| \sum_{i=1}^{m} \sigma_{i} \Phi_{K}(x_{i}) \right\|_{\mathbb{H}_{K}} \right] \\ &= \frac{1}{m} \operatorname{\mathbb{E}} \left[ \sup_{K \in \mathcal{K}} \sqrt{\left\| \left\| \sum_{i=1}^{m} \sigma_{i} \Phi_{K}(x_{i}) \right\|_{\mathbb{H}_{K}} \right]} \\ &= \frac{1}{m} \operatorname{\mathbb{E}} \left[ \sup_{\gamma > 0} \sqrt{\sigma^{\top} \mathbf{K}_{\gamma} \sigma} \right] \\ &= \frac{1}{m} \operatorname{\mathbb{E}} \left[ \sqrt{\sup_{\gamma > 0} \sigma^{\top} \mathbf{K}_{\gamma} \sigma} \right]. \end{aligned}$$

2. Suppose  $\|\mathbf{x}_i - \mathbf{x}_j\| = 1$  for  $i \neq j$ . Compute exactly  $\widehat{\mathfrak{R}}_S(\mathcal{H})$ .

Solution: Given that  $\|\mathbf{x}_i - \mathbf{x}_j\| = 1$  for  $i \neq j$ , the diagonal terms of the kernel matrix are  $\mathbf{K}_{\gamma}^{i,j} = 1$  for i = j and the off-diagonal terms are  $\mathbf{K}_{\gamma}^{i,j} = e^{-\gamma}$  for  $i \neq j$ .

$$\sup_{\gamma>0} \left[ \boldsymbol{\sigma}^{\top} \mathbf{K}_{\gamma} \boldsymbol{\sigma} \right] = \sup_{\gamma>0} \left[ \sum_{i,j} \sigma_{i} \sigma_{j} \mathbf{K}_{\gamma}^{i,j} \right]$$
$$= \sup_{\gamma>0} \left[ m + e^{-\gamma} \sum_{i \neq j} \sigma_{i} \sigma_{j} \right]$$
$$= \sup_{\gamma>0} \left[ \sum_{i,j} \sigma_{i} \sigma_{j} \mathbf{K}_{\gamma}^{i,j} \right]$$
$$= m + \sup_{\gamma>0} e^{-\gamma} \sum_{i \neq j} \sigma_{i} \sigma_{j}$$
$$= m + \sum_{i \neq j} \sigma_{i} \sigma_{j} \mathbf{1}_{\sum_{i \neq j} \sigma_{i} \sigma_{j}>0}.$$

Observe that:

$$m + \sum_{i \neq j} \sigma_i \sigma_j = \sum_{i,j=1}^m \sigma_i \sigma_j = \boldsymbol{\sigma}^\top \mathbf{1} \mathbf{1}^\top \boldsymbol{\sigma} = (\boldsymbol{\sigma}^\top \mathbf{1})^2 = \left[\sum_{i=1}^m \sigma_i\right]^2.$$

It is also known that:

$$\mathbb{E}\left[\left|\sum_{i=1}^{m} \sigma_{i}\right|\right] = \frac{1}{2^{m-1}} \left\lceil \frac{m}{2} \right\rceil \binom{m}{\left\lceil \frac{m}{2} \right\rceil} \le \sqrt{m}.$$
 (Jensen's ineq.)

Thus, we have:

$$\sup_{\gamma>0} \left[ \boldsymbol{\sigma}^{\top} \mathbf{K}_{\gamma} \boldsymbol{\sigma} \right] = \begin{cases} |\sum_{i=1}^{m} \sigma_{i}| & \text{if } \sum_{i \neq j} \sigma_{i} \sigma_{j} > 0; \\ \sqrt{m} & \text{if } \sum_{i \neq j} \sigma_{i} \sigma_{j} < 0; \\ \sqrt{m} & \text{if } \sum_{i \neq j} \sigma_{i} \sigma_{j} = 0. \end{cases}$$

When m is odd, the event  $\sum_{i \neq j} \sigma_i \sigma_j = 0$  cannot occur and the other two events are symmetric, each with probability 1/2. Thus, we have:

$$\widehat{\mathfrak{R}}_{S}(\mathcal{H}) = \frac{1}{2^{m}} \frac{m+1}{2m} \binom{m}{\frac{m+1}{2}} + \frac{1}{2} \frac{1}{\sqrt{m}}.$$

When *m* is even, the event  $\sum_{i \neq j} \sigma_i \sigma_j = 0$  occurs with probability  $\frac{1}{2^m} {m \choose 2}$  and the other two events with equal probability  $p = \frac{1}{2} - \frac{1}{2^{m+1}} {m \choose 2}$ . Thus, we have:

$$\widehat{\mathfrak{R}}_{S}(\mathcal{H}) = \left[\frac{1}{2} - \frac{1}{2^{m+1}} \binom{m}{\frac{m}{2}}\right] \frac{1}{2^{m}} \binom{m}{\frac{m}{2}} + \left[\frac{1}{2} + \frac{1}{2^{m+1}} \binom{m}{\frac{m}{2}}\right] \frac{1}{\sqrt{m}}.$$

We can express the solution in terms of  $\beta_0 \approx \sqrt{\frac{2}{\pi}}$ , where  $\frac{1}{m} \mathbb{E}[|\sum_{i=1}^m \sigma_i|] = \frac{\beta_0}{\sqrt{m}}$ , as follows:

$$\widehat{\mathfrak{R}}_{S}(\mathcal{H}) = \begin{cases} \frac{1}{2} [\beta_{0} + 1] \frac{1}{\sqrt{m}} & \text{if } m \text{ even} \\ \frac{1}{2} [\beta_{0} + 1] \frac{1}{\sqrt{m}} + \frac{1}{2} [\beta_{0} - \beta_{0}^{2}] \frac{1}{m} & \text{otherwise.} \end{cases}$$