

(2021-11-04)

Last time:  $\pi_1(S^n) \cong \emptyset \quad \forall n \geq 2$ . One quick consequence before moving on.

Then (Borsuk-Ulam): Let  $f: S^2 \rightarrow \mathbb{R}^2$ . Then  $\exists x \in S^2$  s.t.  $f(x) = f(-x)$

Pf: (Note: Similar for  $f: S^1 \rightarrow \mathbb{R}$  — let  $g(x) = f(x) - f(-x)$ )

~~This then follows from intermediate value theorem~~ (Claim:  $\exists x \in S^1$  s.t.  $g(x) = 0$ )

Then  $g$  is cts,  $g(x) = -g(-x)$  — so this follows by IUT.

Here, need something ~~better~~ Now, for the sphere  $S^2$ .

Suppose  $\nexists x \in S^2$  s.t.  $f(x) = f(-x)$ . Let  $s(x) = \frac{f(x) - f(-x)}{\|f(x) - f(-x)\|}$  —

this is a cts map  $g: S^2 \rightarrow S^1$ , s.t.  $g(x) = g(-x)$ .

Consider the equator  $\gamma(t) = (\cos(2\pi t), \sin(2\pi t), 0)$ .

and the image  $g \circ \gamma$ . On one hand —  $\gamma \cong \emptyset \Rightarrow g \circ \gamma \cong \emptyset$ .

OTOH,  $\text{VLG } g \circ \gamma(0) = x_0, g \circ \gamma(\frac{1}{2}) = -x_0$ . Consider  $g \circ \gamma: [0, 1] \rightarrow \mathbb{R}$ .

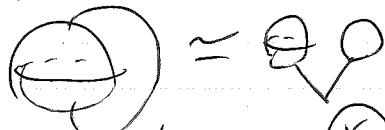
Then  $\gamma \circ \gamma(0) = 0, g \circ \gamma(\frac{1}{2}) = k + \frac{1}{2}$ . Further,  $g \circ \gamma: [\frac{1}{2}, 1] \rightarrow \mathbb{R}$  is a copy of  $g \circ \gamma: [0, \frac{1}{2}] \rightarrow \mathbb{R}$ .

so  $g \circ \gamma(\frac{1}{2} + t) = k + \frac{1}{2} + g \circ \gamma(t)$

$\Rightarrow g \circ \gamma(1) = 2k + 1$  is odd. This contradicts  $g \circ \gamma \cong \emptyset$ .

Today: Combining spaces.

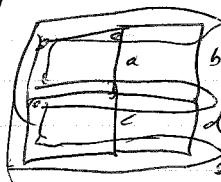
Recall: HTpy equiv:



lots of spaces are htpy equiv to a  bouquet:

Ex: Let  $X$  be a connected finite graph.

Then  $X$  has a  spanning tree  $T$  (a tree that contains every vertex). Spanning trees are maximal



- if you add any edge, it forms a loop.

So if you quotient by the tree

$$\# \text{ of edges} = \# \text{ of vertices} - 1$$

(Same is true for infinite graphs, but have to deal with  $\omega$ 's)

Q:  $\pi_1$  of a bouquet?

Def: Let  $(X, x_0), (Y, y_0)$  be spaces. The wedge sum of  $X$  and  $Y$  is  $X \vee Y = \frac{X+Y}{x_0 \sim y_0}$ . The bouquet of  $X$  and  $Y$  is

$\text{W} = \bigcirc \circ$  If  $X, Y$  are complexes, then these are highly equivalent.  
 We'll work with them because it's a little easier.

If  $X, Y$  are complexes then  $\pi_1(X) * \pi_1(Y)$

Prop:  $\pi_1(X \vee Y)$  is the free product of  $\pi_1(X)$  and  $\pi_1(Y)$ .

Def:  $\pi_1(X) * \pi_1(Y)$  is the group consisting of words of  $S$

A word  $w$  is  $\pi_1(X) * \pi_1(Y)$  if  $w$  is a formal  $w = \prod a_i$  where  $a_i \in \pi_1(X) \cup \pi_1(Y)$ .  
 $i=1 \dots n$

The free group  $\pi_1(X) * \pi_1(Y)$  is the group consisting of words in  $\pi_1(X) \cup \pi_1(Y)$  under the equiv relation of "obvious relations". (GLH)\*

$$wgg'w' \sim w(gg')w'$$

$$whh'w' \sim w(hh')w'$$

$$we_Gw' \sim ww' \sim we_Hw'.$$

with operation concatenation as the group operation

$$\text{So } (a_1 \dots a_n)^{-1} = a_1^{-1} a_2^{-1} \dots a_n^{-1}$$

$$\text{the } e_{G \times H} = e_G = e_H = \emptyset \text{ (empty word)}$$

Special Ex:  $\mathbb{Z} * \mathbb{Z} = F_2$  - free group of rank 2.

$$\langle a \rangle * \langle b \rangle$$

$$F_2 = \{e, a, b, a^{\pm}, b^{\pm}, ab, ba, aa, ab, ab^{\pm}, ba, bb, ba^{\pm}\}$$

Prop: Every element of  $G * H$  has a unique "reduced" form,

$w = \prod a_i$  where  $a_i \in e_G \cup e_H$  and elements of  $G$  and  $H$

alternate

- If  $K$  is a group,  $\alpha: G \rightarrow K$ ,  $\beta: H \rightarrow K$ , then there is a unique homomorphism  $f: G * H \rightarrow K$  st.

$$f(g) = \alpha(g) \quad \forall g \in G$$

$$f(h) = \beta(h) \quad \forall h \in H$$

$$(back) \text{ Let } \pi_1(V) = \pi_1(G) * \pi_1(H)$$

Let  $X, Y$  be path-connected

Prop: Let  $(W)$  be a bouquet of  $(X, x_0)$  and  $(Y, y_0)$ . Then

$$\pi_1(X, x_0) * \pi_1(Y, y_0) \longrightarrow \pi_1(W)$$

(in fact, iso) ~~but not surj~~, but we'll see that later.)

Ex: Let  $\gamma$  be a loop in  $W$  & we want to write  $\gamma$  as a product of loops in  $X$  &  $Y$ . But we need a little work to avoid



Let  $X'$  be a nbhd of left side of

Let  $Y'$  be a nbhd of right side

so that  ~~$X \cap Y$  is contractible~~  $X' \cap Y'$  is contractible  $X'$  retracts to  $X$ ,  $Y'$  retracts to  $Y$

Let  $\gamma: [0, 1] \rightarrow W$ .  
 Let  $U_1 = \gamma(X)$ ,  $U_2 = \gamma(Y)$ . This is an open cover of  $[0, 1]$ .  
 so  $\exists n > 0$  s.t. every  $\frac{1}{n}$ -ball in  $[0, 1]$  lies in  $U_1$  or  $U_2$ .

So cut up  $[0, 1]$ : 

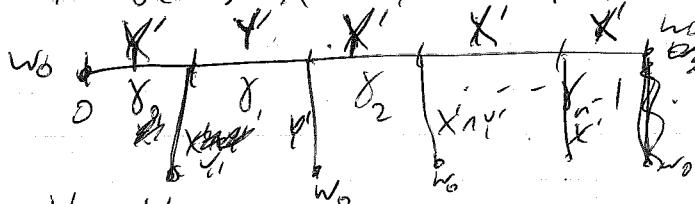
Let  $\gamma_i = \gamma|_{[i/n, (i+1)/n]}$  so  $\gamma = \gamma_0 \circ \gamma_1 \circ \dots \circ \gamma_{n-1}$  where each  $\gamma_i$  is a path in  $X$  or  $Y$ .

$\forall i = 0, \dots, n$ , let  $\lambda_i$  be a path from  $w_0$  to  $\gamma(i/n)$  s.t.

if  $\gamma(i/n) \in X'$ , then  $\lambda_i \subset X'$

if  $\gamma(i/n) \in Y'$ , then  $\lambda_i \subset Y'$

if  $\gamma(i/n) \in X' \cap Y'$ , then  $\lambda_i \subset X' \cap Y'$ .



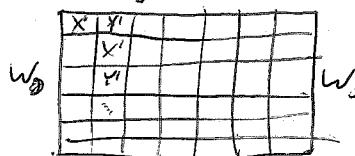
Then  $\lambda_0 \circ \lambda_1 \circ \dots \circ \lambda_{n-1}$  is a loop based at  $w_0$  which lies in  $X'$  or  $Y'$ .  
 And  $\gamma \simeq \lambda_0 \circ \lambda_1 \circ \dots \circ \lambda_{n-1} \circ w_0 w_1 \circ \dots \circ w_{n-1} w_0 \in \pi_1(X) * \pi_1(Y)$ .

What if we want to prove that  $\pi_1(X \vee Y) \cong \pi_1(X) * \pi_1(Y)$ ?

Suppose that  $\gamma = \alpha_1 \circ \dots \circ \alpha_n$ ,  $\alpha_i \in \pi_1(X) * \pi_1(Y)$  and  $\gamma \simeq 0$ .

Need to show that  $\gamma = 0$ , by arcs or loops b-

Idea:



but implementation is quite tricky.

Instead, let's state the general version:

Thm (van Kampen): Let  $X = \bigcup A_\alpha$  is a union of open sets that share a basepoint  $x_0$ . Then  $\pi_1(X) \cong \pi_1(A_1) * \pi_1(A_2) * \dots$

~~If each intersection  $A_\alpha \cap A_\beta$  is path-connected,~~

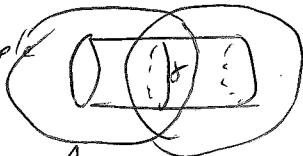
then let  $\Phi: \ast \pi_1(A_\alpha) \rightarrow \pi_1(X)$  be the map extending the maps induced by  $i_\alpha$  that agrees with  $i_\alpha$  on  $\pi_1(A_\alpha)$  where  $i_\alpha: A_\alpha \hookrightarrow X$ .

Let's say for  $[f] \in \pi_1(A_\alpha)$ , with  $[f]_\alpha = \underline{\underline{f}}([f])$ .

Then: ① If  $A_\alpha \cap A_\beta$  is path-connected  $\forall \alpha, \beta \neq$  then  $\Phi$  is injective.

( $\exists$ ) Typically not injective & Trivial example

$$[\gamma] \in \pi_1(A, \gamma A_2)$$



$A_{\alpha 1} \quad A_{\alpha 2}$

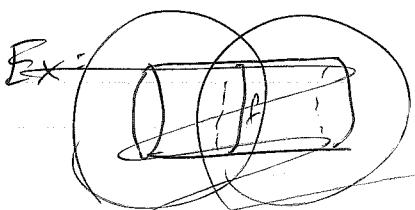
So  $\pi_1(A_1) * \pi_1(A_2)$  contains two copies of  $\gamma - [\gamma] \in \pi_1(A_1)$ ,  
 $[\gamma]_2 \in \pi_1(A_2)$ . And  $\Phi([\gamma]) = (i_1)_*([\gamma]) = [\gamma]_1$ ,  
 $\Phi([\gamma]_2) = (i_2)_*([\gamma]_2) = [\gamma]_2$

- same curve. That is,  $[\gamma], [\gamma]_2 \in \ker \Phi$ .  
(and since the kernel is normal,  $w[\gamma][\gamma]_2^{-1}w^{-1} \in \ker \Phi$  b/w.)

Second part of this is that these are the only things in the kernel.

② If  $A_\alpha \cap A_B \cap A_\beta$  is path-connected for all  $\alpha, \beta, \delta$ , then

$$\begin{aligned}\ker \Phi &= \langle w[f]_a [f]_B^{-1} w^{-1} \mid f \in \pi_1(A_\alpha \cap A_\beta \cap A_\delta) \rangle \\ &= \langle \langle [f]_\alpha [f]_\beta^{-1} \mid f \in \pi_1(A_\alpha \cap A_\beta) \rangle \rangle_{\alpha, \beta}^{w \in \pi_1(A_\delta)}\end{aligned}$$



And therefore,  $\pi_1(X) \cong \pi_1(A_\alpha) / \ker \Phi$ .

$$\begin{aligned}\pi_1(A) &= \langle a \rangle & \pi_1(A \cap B) &= \langle \rangle \\ \pi_1(B) &= \langle b \rangle & \pi_1(A \cap C) &= \langle \rangle \\ \pi_1(A) * \pi_1(B) &= \langle a, b \rangle\end{aligned}$$

$$X = A \cup B$$

$$\pi_1(X) \cong \Phi(a) * \Phi(b)$$

For each element of  $\pi_1(A \cap B)$ ,  $\ker \Phi \subset \text{abs} \rightsquigarrow$

Ex:  we have elements

$$w = \langle x \rangle \quad \Phi: \pi_1(X) * \pi_1(Y) \rightarrow \pi_1(W)$$

$$\ker \Phi = \langle \langle [f]_X [f]_Y^{-1} \mid f \in \pi_1(X \cap Y) \rangle \rangle = \langle \rangle$$

Non-trivial is harder.

Def: Let  $X, Y$  be a space,  $f: X \rightarrow Y$  a map. The mapping cylinder  $C_f$  is  $C_f = Y \cup (X \times [0, 1])$   $(x, 0) \sim f(x)$ .

Then  $C_f$  deformation retracts to  $Y$ .

$$\text{Let } f: (S^1, *) \rightarrow (X, x_0) \quad S^1 \rightsquigarrow \quad Z = C_f \cup C_g \quad (x, 1) \sim (x, 0)$$

$$\pi_1(Z)?$$

$$g: (S^1, *) \rightarrow (Y, y_0)$$

$$\begin{aligned}\pi_1(X, x_0) &\cong \pi_1(X) \\ \pi_1(Y, y_0) &\cong \pi_1(Y)\end{aligned}$$

$$\pi_1(X, z_0) \cong \pi_1(X)$$

$$\pi_1(Y, z_0) \cong \pi_1(Y)$$

$X \cap Y$  is path-connected, so this applies and

$$\cong \pi_1(X) * \pi_1(Y) / \ker \Phi$$

$$\cong \pi_1(X) * \pi_1(Y) / \langle \langle f_g^{-1} \rangle \rangle$$

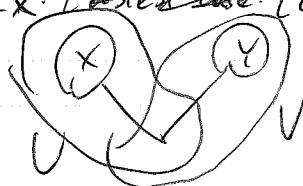
Last time:

Thm (Van Kampen): Let  $X = \bigcup A_\alpha$  be a union of open sets that share a base point  $x_0$ . Let  $\Phi: \prod_{\alpha} \pi_1(A_\alpha, x_0) \rightarrow \pi_1(X, x_0)$  be the canonical map.

① If  $A_\alpha \cap A_\beta$  is path-connected  $\forall \alpha, \beta$ , then  $\Phi$  is surjective.

② If  $A_\alpha \cap A_\beta \cap A_\gamma$  is path-connected  $\forall \alpha, \beta, \gamma$ , then  $\ker \Phi = \langle [f]_\alpha [f]_\beta^{-1} \mid f \in \pi_1(A_\alpha \cap A_\beta, x_0) \rangle$ .

Ex: Easiest case: let  $X = U \cup V$  where  $U \cap V$  is path-connected.



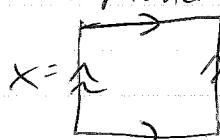
Simplest:  $\pi_1(U \cap V) = 0$ .

$X = U \cup V$ , intersections are path-connected.

$$\Rightarrow \pi_1(X) \cong \pi_1(U) * \pi_1(V)$$

$\cong \pi_1(X) * \pi_1(Y)$

More complicated:  $\pi_1(U \cap V) \cong \mathbb{Z}$ .



$U \cap V = \text{circle}$

And left because



$U$  retracts to



$$\text{Actually, let's be careful about loops!} \quad \pi_1(U) = \langle a, b \rangle$$

$$\pi_1(V) = 0$$

$$\pi_1(X) = \langle a, b \rangle * 0$$

$V$  retracts to,  $V = \text{circle}$   $\pi_1(V) = \langle aba^{-1}b^{-1} \rangle$

$$[\gamma]_U [\gamma]_V = aba^{-1}b^{-1} = \langle a, b \rangle$$

This guarantees essentially  $\pi_1(X) = \langle a, b \rangle$

$$\text{Written } \pi_1(X) \cong \langle a, b \mid aba^{-1}b^{-1} = 1 \rangle$$

$$\text{or } \pi_1(X) \cong \langle a, b \mid ab = ba \rangle$$

$$\text{i.e. } \pi_1(X) \cong \mathbb{Z}^2.$$

Likewise:

$$\pi_1(X) = \langle a, b, c \mid aba^{-1}b^{-1}c^{-1} = 1 \rangle$$

Ex:



$$\pi_1(X) = \langle a, b, c \mid aba^{-1}b^{-1}c^{-1} = 1 \rangle$$

A couple ways: (1) or:

$$\pi_1(U) = \langle a, b \rangle \quad \pi_1(V) = \langle c, d \rangle$$

$$U \cap V = \text{circle} \quad \pi_1(U \cap V) = \langle aba^{-1}b^{-1}c^{-1} \rangle$$

$$[\gamma]_v = cd c^{-1} d^{-1}, \pi_v(k) = \frac{\pi_v(U) * \pi_v(V)}{\langle [\gamma]_v, [\gamma_v]^{-1} \rangle}$$

$$= \langle a, b, c, d \mid aba^{-1}b^{-1} = cd c^{-1} d^{-1} \rangle$$

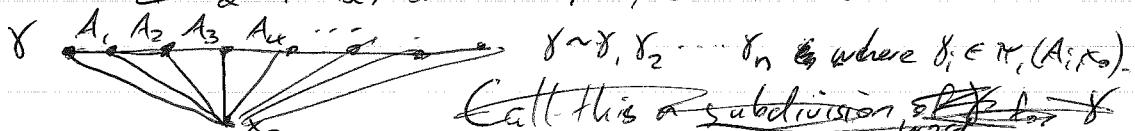
More This is an example of an amalgamated free product:

$$\text{if } A \hookrightarrow G \quad G *_A H = \frac{G * H}{i(a) * j(a)}$$

$$\hookrightarrow H_a \quad = G * H \frac{}{\langle i(a)j(a)^{-1} \rangle}$$

Pf:

Previously if  $X = \bigcup A_\alpha$  and  $A_\alpha \cap A_\beta$  is path-connected, then  $\Phi = \pi_1(A_\alpha, x_0) \rightarrow \pi_1(X, x_0)$  is surjective.



Call this a subdivision word

② Need to translate homotopies. Suppose  $\gamma \in \pi_1$ .

$$\text{Suppose } w = [f] \dots [f_n]_{\alpha_n} \sim$$

$$\text{Also } w' = [f'_1] \dots [f'_n]_{\alpha'_n}.$$

(Claim:  $w'$  can be transformed to  $w'$  do a series of ops:

(A)  $v[\gamma]_\alpha [\gamma']_\alpha v' \leftrightarrow v[\gamma\gamma']_\alpha v'$

(B)  $v[e]_\alpha v \leftrightarrow v'$

(C)  $v[\gamma]_\alpha v' \leftrightarrow v[\gamma]_\beta v'$  if  $\beta \in \pi_1(A_\alpha \cap A_\beta)$ .

Pf: Take top 3 bits

Ops

Subdivide into cells

at: Image of each cell lies in some  $A_\alpha$

- no more than 3 edges at a vertex

- The vertices include the original verts.

Lemma: If  $\gamma$  is a curve in  $X \in \pi_1(X)$ , then  $\gamma$  and  $\gamma'$  are homotopic if and only if  $w, w'$  are two different subdivision words for  $\gamma$  and  $\gamma'$  respectively.

Pf:

(A)  $\gamma \rightarrow \gamma'$

(B)  $\gamma \rightarrow \gamma'$

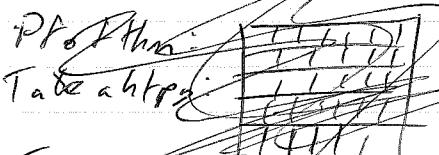
Take a common subdivision

Ops: - change tails.

Ops: - add vertex

- delete vertex

- change labeling



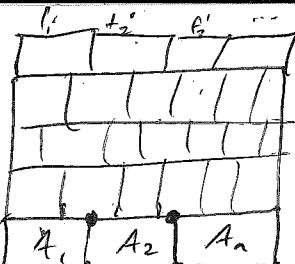
Subdivide so that - change labeling

each cell lies in some  $A_\alpha$

- no more than 3 edges at a vertex.

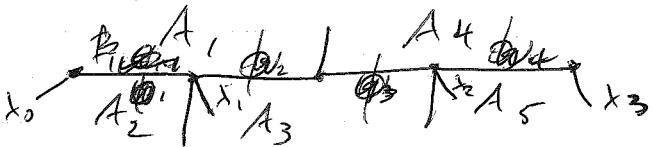
For

PF: Subdivide a hexagon:



- each cell lies in some  $A_\alpha$ .
- no more than 3 edges at each vertex.
- The vertices include the original vertices

Add a tail  $f_1, f_2, \dots, f_{n+2}$  to each vertex so that  $\lambda_v = \lambda_v f_1 f_2 \dots f_{n+2} \in \pi_1(A_\alpha \cap A_\beta)$ .  
 for each edge we have a loop  $\gamma_e = \lambda_v, g_e$  word.  
 Then every edge path  $\gamma_{e_1 e_2 \dots e_m}$  can be read as a word.  
 In fact, typically in multiple equivalent ways:  
 Top line is  $w^*$ , Bottom is  $w$ .



$$[(x_0 x_1) x_2] x_3 [(x_4 x_5) x_6] x_5 \sim [x_0, x_1, x_2]$$

but these are equivalent by move (B).

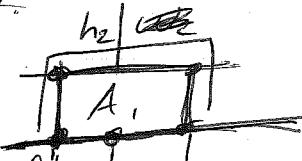
Top line is  $w^*$ , bottom is  $w$ .

Add a tail for each chunk,

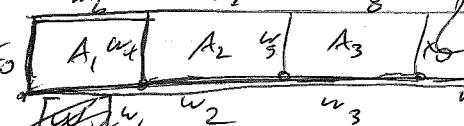
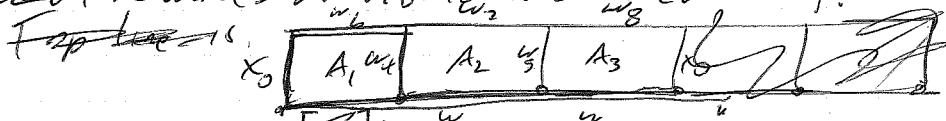
For each 2-cell,

the word reading around circumference is

is the ~~identical~~ equal to the identity by move (A).



So: The words on horizontal lines are all equivalent.



$$w_1 w_2 w_3 \sim w_6 w_4 w_2 w_3 \sim w_6 w_5 w_3 \sim w_6 w_2 w_4$$

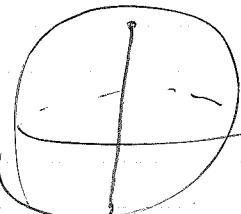
top line is  $w^*$ , bottom is  $w$ .

(break)  
Overflow.

Fundamental groups of complexes:

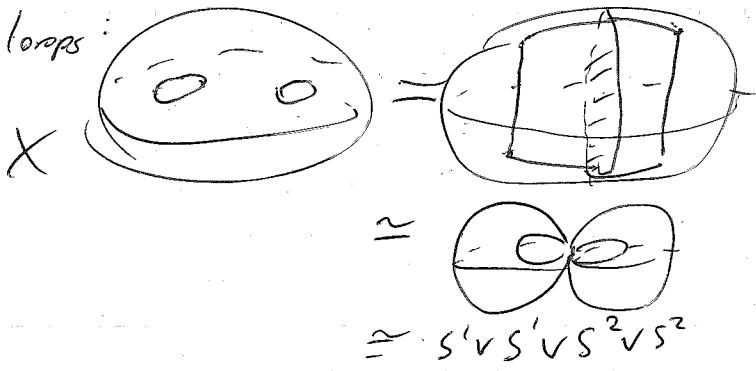
Couple examples w/ knot / link complements:

$$\mathbb{R}P^3 \setminus S^1$$



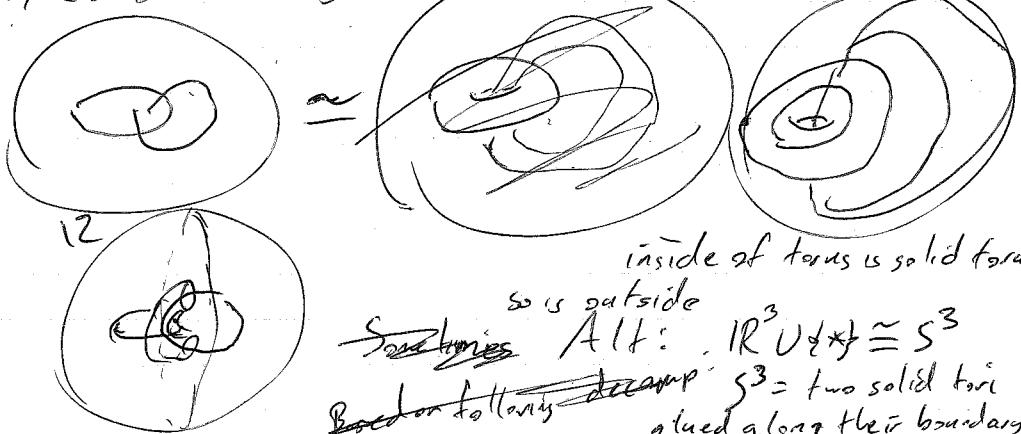
$$\text{By KT, } \pi_1(S^1 \vee S^2) \cong \mathbb{Z}.$$

Say we have two loops:



$$\pi_1(X) \cong \mathbb{Z} * \mathbb{Z}$$

But:



Maybe:

~~Stringy Alt:~~  $\mathbb{R}^3 \setminus \{*\} \cong S^3$

~~Based on torus decomps~~  $S^3 = \text{two solid tori glued along their boundary.}$

~~One more~~ Torus knots: Recall:  $\pi_1(T) \cong \mathbb{Z} * \mathbb{Z}$

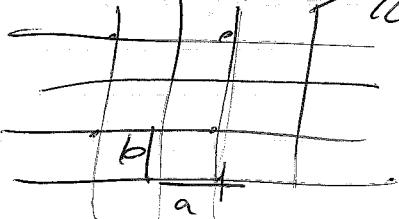
Embed  $T$  in  $\mathbb{R}^3$ :

relatively prime. Then:

Then  $\text{Rep}(T)$  can be represented by a simple curve on  $T$  ( $\mathbb{C}$ )

$m, n \in \mathbb{Z}$  are relatively prime. w.t.  $|w| \leq \frac{1}{2}$

We can write  $T = \mathbb{R}^2 / \mathbb{Z}^2$



$\forall m, n \in \mathbb{Z}$

$\gamma: [0, 1] \rightarrow \mathbb{R}^2 / \mathbb{Z}^2$

$\gamma(t) = (mt, nt)$

is a closed curve s.t.

$\gamma \sim a^m b^n$ . If  $\text{gcd}(m, n) = 1$ ,

this curve is simple (no self-intersections).

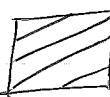
Conversely, if  $\gamma: S^1 \rightarrow \mathbb{R}^2 / \mathbb{Z}^2$  is a simple closed curve, then there is  $a^m b^n$  where  $\text{gcd}(m, n) = 1$ .

Embed  $T$  in  $\mathbb{R}^3$  in standard way:

The  $(m, n)$ -torus knot is  $\mathbb{R}^3 \setminus \gamma$ , where  $\gamma \subset T$

where  $\gamma$  is isotopic to  $a^m b^n$ .

$\pi_1(\mathbb{R}^3 \setminus K_{m,n})$



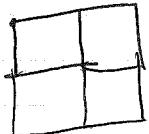
2021-11-18

## Fundamental groups of complexes:

1-dim'l: Let  $X$  be a graph. It has a spanning tree  $T \subset X$  which is contractible. Then  $X \cong X/T \cong \mathbb{B}$ : one loop per edge.

$\pi_1(X) \cong \langle e_1, \dots, e_k \rangle$  where  $k = \# \text{ of edges in } X - T$   
 free group with one generator for each edge not in  $T$ ,  
 $\Rightarrow \text{rank}(\pi_1(X)) = E - V + 1$ .

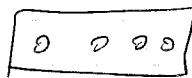
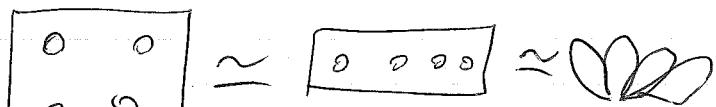
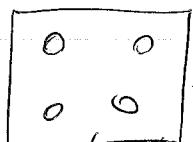
Alt:



~~contract~~

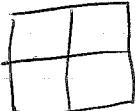
if  $X$  is planar, ~~then~~

$\sim$



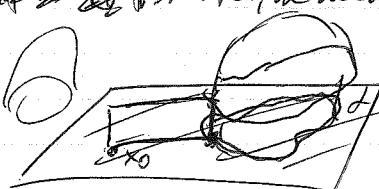
- one generator for each interior face.

2-dim'l:



Idea: pop the 2-cells, remove a point in each cell, rest retracts to 1-skeleton, glue on a patch that covers the pt.

But ~~for 2kt~~, we need base points. These base points are all different



These don't share a base point. So we can fix this, add cells, change the set, but it's easier to change the space.

Let  $Y = X \cup D$

Then let  $\gamma$  be a path from  $\text{ex}(y_0)$  to  $x_0$ .

(glue a rectangle) Let  $X' = X \cup B \cup Z$  where



Then let  $\alpha: \partial D \rightarrow Z_1$ . Then  $X \cup D \cong X' \cup D$

For each  $x$ , glue a step rectangle to  $X$  along  $\gamma_x$ . One edge glues to  $\gamma_x$ , one to a segment in  $D_x$ .

Call this space  $Z$ . For each  $x$ , let  $c_x$  be a point in  $D_x$ .

Let  $T = Z \setminus \{c_x \mid x \in X\}$  puncturing every 2-cell

Then  $\pi_1(S) = \pi_1(X)$

$\pi_1(T) = \emptyset$

$\pi_1(S \cap T) = \langle g_{\alpha}, x \in \mathbb{Z} \rangle$

$$\text{And so: } \check{\pi}_1(KF) : \check{\pi}_1(Z) \cong \check{\pi}_1(X) * \check{\pi}_1(\partial)$$

$$\cong \langle [g_1] \circ [g_2]^{-1} \rangle$$

$$\cong \langle \gamma_2 \circ \gamma_2^{-1} \rangle$$

- And we can compute fundamental groups of arbitrary 2-complexes.  
 (good to try this yourself, will put a problem on pset)

Ex: If  $X$  is a complex,  $X^{(2)}$  its 2-skel, then  $\check{\pi}_1(X) \cong \check{\pi}_1(X^{(2)})$ .

$$RP^n = S^n / R_2 \sim x \sim -x \quad \forall x \in S^n. \text{ Spheres } S^n / R_2$$

Cellulate by starting w/ a cellulation of  $S^n$  - with action  $x \mapsto -x$

$$S^2 \text{ w/ 2 cells in each dim from 0 to } n.$$

So Then  $R_2$  ~~acts~~ -action sends cell to cells.  
 And  $S^n / R_2$  has one cell in each dimension.

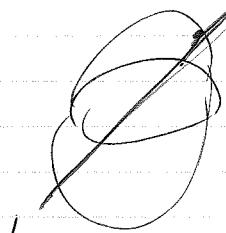
$$RP^0 \text{ } RP^1 \text{ } RP^2. \quad \text{For } n \geq 2, \check{\pi}_1(RP^n) \cong \check{\pi}_1(RP^2)$$

$$\cong \langle \text{ (cell)} \rangle \cong \langle \text{ (cell)}^2 \rangle \cong R_2.$$

(And the same calc works with any cellulation - fewer cells make it easier.)

Alternatively, for a geometric argument:

$$RP^n \cong \overline{R^n \setminus \{0\}} / R_2 = \{ \text{lines through the origin in } R^n \}$$



Pick a "basepoint"  $L_0$ . Consider a loop based at  $L_0$ .  
 This is a path that starts and ends at  $L_0$ .  
 There are two possibilities:

- going around the loop ~~flips~~ preserves orientation of the line
- the loop ~~flips~~ orientation

Then:- Loops that preserve orientation are null-homotopic  
 (they correspond to a loop on the sphere)

- Loops that flip orientation are not null-homotopic  
 (every int pt loop also flips too).

- Compositions act Concatenations of loops act like ~~groups~~ like classes ~~preserves~~ a group. (flip, then flip = preserve, etc).

In fact, there ~~is~~ just  $\check{\pi}_1(RP^1)$ .

And we can describe this in terms of lifting:

$$\text{Prev: } \begin{array}{ccc} \widetilde{Y} & \xrightarrow{\quad} & \mathbb{R} \\ & \downarrow p & \\ [0,1] & \xrightarrow{\quad} & S^1 \end{array} \quad \begin{array}{ccc} \widetilde{Y} & \xrightarrow{\quad} & S^n = \text{oriented lines} \\ & \downarrow p & \\ [0,1] & \xrightarrow{\quad} & \mathbb{RP}^n = \text{unoriented lines} \end{array}$$

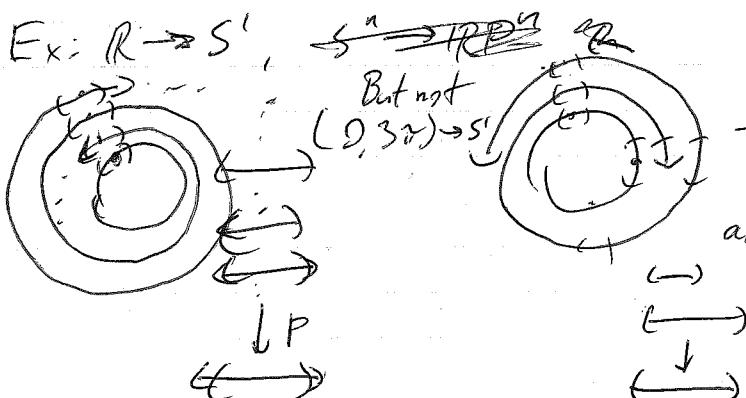
(paths lift, ~~the~~ homotopy classes correspond to endpoint lifting)

Can we generalize?

Covering spaces: A covering space  $\tilde{X}$  of  $X$  is a space  $\tilde{X}$  w/ a map  $p: \tilde{X} \rightarrow X$  s.t.  $\forall x \in X$ ,  $\exists U \ni x$  s.t.  $p^{-1}(U)$  is a disjoint union of copies of  $U$ , ~~s.t.~~ each copy of  $U$  is sent homeomorphically to  $U$  by  $p$ .

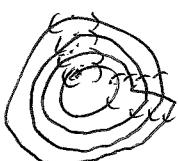
So,  $\tilde{p}^{-1}(U) \cong U \times E$  where  $E$  is discrete.

~~These are called sheets of  $\tilde{X}$~~



Ex:  $S^1 \rightarrow S^1$  ( $S^1 = \{z \in \mathbb{C} \mid |z| = 1\}$ )

$$z \mapsto z^n$$



Ex:  $S^n \rightarrow \mathbb{RP}^n$  If we take a ~~line~~ a collection of nearby lines, then Every <sup>unoriented</sup> line has two oriented lines. If we take a line  $L$ , nbhd  $U \ni L$ , ~~then~~  $p^{-1}(U)$  is two copies of  $U$ , one oriented one way, one oriented the other way.

Key fact then:

Lifting: If  $p: (\tilde{X}, \tilde{x}_0) \rightarrow (X, x_0)$  is a cover, and  $f: [0,1]^n \rightarrow X$  is a map, s.t.  $f(0) = x_0$ ,

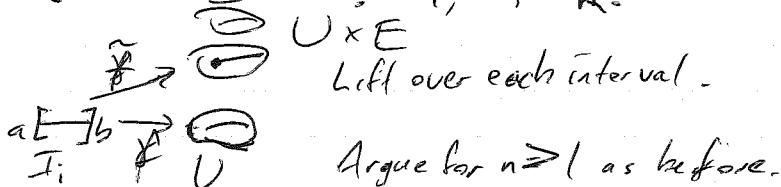
$\tilde{f} : [0, 1]^n \rightarrow X$   
 there is a unique lift  $\tilde{f}$  s.t.  $\tilde{f}(0) = x_0$ .

Pf: (a=1) Lift short segments. Subdivide and lift.

$\forall x \in X, \exists U_x$  s.t.  $\rho^{-1}(U_x) \cong U_x \times E_x$ ,  $E_x$  discrete,

$p(u, e) = u$ .  
 Then  $f^{-1}(U_x)$  is an open cover of  $[0, 1]$ .  $\exists \varepsilon > 0$  s.t. each  $f([a, a+\varepsilon])$  is contained in some  $U_x$ .

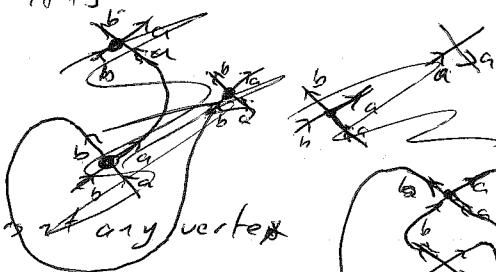
Subdivide  $[0, 1]$  into  $\varepsilon$ -intervals,  $I_1, \dots, I_k$ .



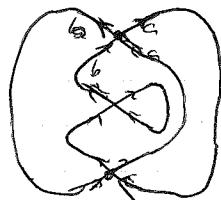
This is why  
 $(0, 3\pi) \rightarrow S^1$   
 doesn't work

And we can use this point of view to construct  
 more complex lifts.

$$X = \text{a figure-eight graph}$$



a and b have lifts starting at any vertex.



Follow those lifts, end at a vertex.

- Any graph with edges labeled a and b, with one a-edge in and out of each vertex, one b-edge in and out of each type in and out of any vertex is a cover of X.

Then (Fundamental Theorem of Covering Spaces): If  $X$  is a connected complex, manifold,  
 Then there is a bijection between {connected based covers} of  $(X, x_0)$  and {subgroups of  $\pi_1(X, x_0)$ }.

$$(X, x_0) \xrightarrow{\quad} \rho \ast (\pi_1(X, x_0)).$$

Ex: cube, face group.