

1. Show how to encode two classical bits into one qubit such that any one bit can be recovered correctly with probability greater than 85%.
2. We have seen in class that Toffoli gates are universal for classical reversible computation. Prove that no set of two-bit and one-bit gates is universal for classical reversible computation.
3. Let $f : \{0, 1\}^n \rightarrow \{0, 1\}^m$, $g : \{0, 1\}^m \rightarrow \{0, 1\}^n$ be such that g is a left-inverse of f (i.e., $g(f(x)) = x$ for all $x \in \{0, 1\}^n$). Assume that both functions can be computed by a polynomial size classical circuit. Show that there exists a polynomial size classical reversible circuit (and hence also quantum circuit) that maps $|x, 0\rangle$ to $|f(x), 0\rangle$. Would you expect this to be possible without the assumption that g has a polynomial size circuit?
4. Describe a quantum algorithm that solves the following problem. Given a function $f : \mathbb{Z}_2^n \rightarrow \{0, 1\}^m$ that satisfies $f(x) = f(y) \Leftrightarrow x - y \in H$ for some subgroup H of \mathbb{Z}_2^n , find H .
5. For any function $f : \{0, 1\}^n \rightarrow \{0, 1\}$ we define U_f as the unitary transformation mapping $|x, y\rangle$ to $|x, y + f(x)\rangle$ for each $x \in \{0, 1\}^n$ and $y \in \{0, 1\}$. Also define S_f as the unitary transformation mapping $|x\rangle$ to $(-1)^{f(x)}|x\rangle$ for each $x \in \{0, 1\}^n$. Show how to obtain S_f from U_f (using an auxiliary qubit). Can you obtain U_f from S_f ?
6. (a) Let $|u\rangle, |v\rangle$ be two states on n qubits each. Consider the circuit below, which uses a controlled swap gate. Find the probability of measuring $|0\rangle$ as a function of $|\langle u|v\rangle|$. What does this quantity correspond to?

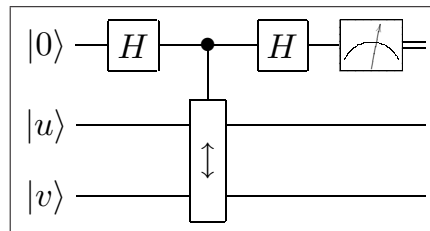


Figure 1: The swap test

- (b) Now assume there exists a quantum circuit U that transforms $|0\rangle$ to $|u\rangle$ and a quantum circuit V that transforms $|0\rangle$ to $|v\rangle$. Show how to generate the state $(|0\rangle|u\rangle + |1\rangle|v\rangle)/\sqrt{2}$ using U and V . Then, assume we apply H on the first qubit and measure it. Find the probability of measuring $|0\rangle$ as a function of $|\langle u|v\rangle|$.
7. Here we develop parts of the very useful *phase estimation* technique, due to Kitaev. Let U be a unitary transformation on n qubits and let $|v\rangle$ be an eigenvector of U with eigenvalue λ .
 - (a) Show that $|\lambda| = 1$, i.e., there exists some $\theta \in [0, 2\pi)$ such that $\lambda = e^{i\theta}$.
 - (b) Based on the circuit shown in Figure 2, describe how to estimate θ to within some additive error ε (with confidence, say, 90%). You can assume that you have a way to generate the state $|v\rangle$. How many operations are needed (roughly) as a function of ε ?

- (c) Show that you can do the same even if you are given only one copy of $|v\rangle$ (and you are unable to generate more yourself).

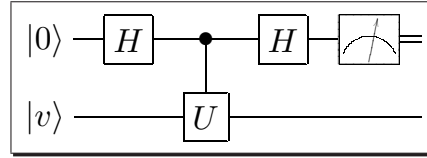


Figure 2: Phase estimation

8. (a) For $b \in \{0, 1\}$ define $|\psi_b\rangle$ as the two-qubit state $\frac{1}{\sqrt{2}}(|00\rangle + (-1)^b|11\rangle)$. Alice and Bob share the state $|\psi_b\rangle$ for some unknown b . Their goal is to determine b . Unfortunately, they are unable to communicate with each other. Convince yourself that Alice (or Bob) cannot determine b alone (no rigorous proof of this is required). Now, assume each of them is allowed to send one classical envelope to a common friend Charlie. Find a protocol that allows Charlie to determine b from the two envelopes he receives.
- (b) For each $k, l \in \{0, 1\}^{2n}$, $k \neq l$, $b \in \{0, 1\}$, define $|\psi_{k,l,b}\rangle$ as the state $\frac{1}{\sqrt{2}}(|k\rangle + (-1)^b|l\rangle)$ on $2n$ qubits. Alice and Bob share the state $|\psi_{k,l,b}\rangle$ for some unknown k, l, b (i.e., each has n qubits). Each of them can send one classical envelope to Charlie, who happens to know k and l but not b . Upon receiving the two envelopes, Charlie is asked to determine the bit b . Find a protocol that allows them to achieve this goal.