The Dehn function of $SL(n; \mathbb{Z})$

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Dehn functions of $SL(n; \mathbb{Z})$

- ▶ *SL*(2; ℤ) is virtually free linear Dehn function
- ► *SL*(3; ℤ) exponential Dehn function (Thurston-Epstein)
- ▶ $SL(4; \mathbb{Z})$ conjectured to be quadratic (Thurston)
- $SL(5+;\mathbb{Z})$ quadratic (Y)

Sol₃: Exponential distortion

$$\mathsf{Sol}_3 = \begin{pmatrix} e^t & 0 & x \\ 0 & e^{-t} & y \\ 0 & 0 & 1 \end{pmatrix} = \mathbb{R} \ltimes \mathbb{R}^2 = \langle t \rangle \ltimes \langle x, y \rangle$$

 \mathbb{R}^2 is exponentially distorted:



Sol₃: Exponential Dehn function



Length:
$$\sim s$$

Area: $\sim e^{2s}$
 $\sim \delta_{Sol_3}(n) \sim e^n$
BUT

$$\mathsf{Sol}_5 = egin{pmatrix} e^{t_1} & & x \ & e^{t_2} & & y \ & & e^{t_3} & z \ & & & 1 \end{pmatrix},$$

 $t_1 + t_2 + t_3 = 0$ has quadratic Dehn function.

The Steinberg presentation

Generators:

$$e_{ij}(x) = egin{pmatrix} 1 & x & \ & 1 & \ & & \ & & \ddots & \ & & & 1 \end{pmatrix} = [e_{ij}(1)]^x = e_{ij}^x$$

Relations:

$$[e_{ij}, e_{k\ell}] = I \text{ if } i \neq \ell, j \neq k$$
$$[e_{ij}, e_{jk}] = e_{ik} \text{ if } i \neq k$$
$$(e_{12}e_{21}^{-1}e_{12})^4 = I$$

The Dehn function of $SL(2; \mathbb{Z})$

$SL(2; \mathbb{Z})$ is virtually free \Rightarrow linear Dehn function. $SL(2; \mathbb{Z})$ is virtually free \Rightarrow powers of generators are far from *I*:

$$d_{SL(2;\mathbb{Z})}(I, e_{ij}(x)) \sim x,$$

BUT

$$d_{SL(3;\mathbb{Z})}(I,e_{ij}(x)) \sim \log |x|.$$

Why?

Distortion in $SL(3; \mathbb{Z})$

$$F = \begin{pmatrix} 2 & 1 \\ 1 & 1 \end{pmatrix}$$
.

Then if

Let

$$H = \left\{ \begin{pmatrix} F^t & x \\ 0 & 0 & 1 \end{pmatrix} \middle| x, y, t \in \mathbb{Z} \right\} \subset SL(3; \mathbb{Z}),$$

then

$$H \cong \begin{pmatrix} \lambda^t & 0 & x \\ 0 & \lambda^{-t} & y \\ 0 & 0 & 1 \end{pmatrix} \subset \mathsf{Sol}_3$$

In particular,

$$\delta_{H}(n) \sim e^{n} \rightsquigarrow \delta_{SL(3;\mathbb{Z})}(n) \sim e^{n}$$

Why should $\delta_{SL(4;\mathbb{Z})}$ be quadratic?

- 1. Part of a larger conjecture by Gromov: Lattices in high-rank symmetric spaces should have polynomial Dehn functions, because there are enough flats.
- 2. By the analogy between $SL(n; \mathbb{Z})$ and Sol_{2n-3} .

 $SL(n;\mathbb{Z})$ and Sol_{2n-3}

$$\mathsf{Sol}_{2n-3} = \left\{ \begin{pmatrix} a_1 & & & z_1 \\ & a_2 & & & z_2 \\ & & \ddots & & \vdots \\ & & & a_{n-1} & z_{n-1} \\ & & & & 1 \end{pmatrix} \middle| \prod a_i = 1 \right\} \subset SL(n; \mathbb{R}).$$

There is a lattice $H \subset \text{Sol}_{2n-3}$ such that $H \subset SL(n; \mathbb{Z})$.

SL(3; ℤ) :

- Contains copies of Sol₃.
- ► There are words *ê*_{ij}(x) of length ~ log |x| which represent e^x_{ij}. (exponential distortion)
- $\delta([\hat{e}_{13}(x), \hat{e}_{23}(x)]) \sim x^2$ (exponential Dehn function)

 $SL(4; \mathbb{Z}):$

- Contains copies of Sol₅.
- There are words $\widehat{e}_{ij}(x)$ of length $\sim \log |x|$ which represent e_{ij}^x .

δ([ê₁₄(x), ê₂₄(x)]) ~ (log |x|)² (quadratic fillings for some words)

 $SL(5;\mathbb{Z}):$

- Contains overlapping solvable groups with quadratic Dehn functions.
- There are words $\widehat{e}_{ij}(x)$ of length $\sim \log |x|$ which represent e_{ij}^x .
- ▶ Different choices of *ê*_{ij}(x) can be connected by quadratic-area homotopies.
- "Simple" words have quadratic fillings.

"Simple" words have quadratic fillings

Theorem (Y)

In $SL(5,\mathbb{Z})$, words of the forms

- $[\widehat{e}_{ij}(x), \widehat{e}_{k\ell}(y)]$ for $i \neq \ell, j \neq k$,
- $[\widehat{e}_{ij}(x), \widehat{e}_{jk}(y)]\widehat{e}_{ik}(xy)^{-1}$ for $i \neq k$,
- $\triangleright \ \widehat{e}_{ij}(x)\widehat{e}_{ij}(y)\widehat{e}_{ij}(x+y)^{-1}.$

have quadratic (~ $(\log |x| + \log |y|)^2)$ filling areas.

Breaking words into simple words

Let $\mathcal{E} = SL(n; \mathbb{R})/SO(n)$ (symmetric space).

- 1. Let w be a word in $SL(n; \mathbb{Z})$; this is also a curve in \mathcal{E} .
- 2. Fill w with a disc and triangulate that disc by triangles of side length $\sim 1.$
- 3. Map the triangulation into $SL(n; \mathbb{Z})$, one dimension at a time.
 - Dim. 0 It suffices to have a map $\mathcal{E} \to SL(n; \mathbb{Z})$: fix a fundamental domain S and map any vertex contained in gS to g.
 - Dim. 1 If x, y are the ends of an edge, then $d(x, y) \leq 1$. We can see where nearby points go by looking at $SL(n; \mathbb{Z}) \setminus \mathcal{E}$.
 - Dim. 2 We can replace the edges of the triangulation by "simple" words, then use the theorem above to fill each one with a quadratic filling.

 n^2 triangles, each filled with area n^2 , leads to a quartic filling inequality