Filling multiples of embedded curves and quantifying nonorientability

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How is FA(T) related to FA(2T)?

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▶ n = 3: If T is a curve in  $\mathbb{R}^3$ , then FA(2T) = 2FA(T). (Federer, 1974)

• n = 4: There is a curve  $T \in \mathbb{R}^4$  such that

 $FA(2T) \leq 1.52 FA(T)$ 

(L. C. Young, 1963)

Let K be a Klein bottle



Let K be a Klein bottle and let T be the sum of 2k + 1 loops in alternating directions.





 T can be filled with k bands and one extra disc D

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$$FA(T) \approx \frac{\text{area } K}{2} + \text{area } D$$





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- $FA(T) \approx \frac{\text{area } K}{2} + \text{area } D$
- 2T can be filled with 2k + 1 bands
- ► FA(2T) ≈ area K— less than 2 FA(T) by 2 area D!

### The main theorem

#### Q: Is FA(2T) bounded below by a function of FA(T)?

- Q: Is FA(2T) bounded below by a function of FA(T)? Theorem (Y.)
- Yes! For any d, n, there is a c > 0 such that if T is a d-cycle in  $\mathbb{R}^n$ , then  $FA(2T) \ge c FA(T)$ .

### Proving the theorem in dimension 0

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pseudo-orientation

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so  $B \mod 2$  is a cycle. If P is an integral cycle such that  $B \equiv P \pmod{2}$  (a *pseudo-orientation* of B), then

$$B + P \equiv 0 \pmod{2}$$
$$\frac{B + P}{2} = \frac{2T + 0}{2} = T.$$



## The Klein bottle, again



### Nonorientability

If A is a mod-2 cycle, define the *nonorientability* of A by

 $NO(A) = \inf\{\max P \mid P \text{ is an integral cycle and } P \equiv A \pmod{2}\}$ 

This measures how hard it is to "lift" A to an integral cycle.

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This measures how hard it is to "lift" A to an integral cycle. If  $\partial B = 2T$ , then

$$\mathsf{FV}(T) \leq \frac{\mathsf{mass}\,B + \mathsf{NO}(B \bmod 2)}{2}$$

So, to prove that  $FV(T) \lesssim FV(2T)$ , it suffices to show: Theorem

If A is a mod-2 d-cycle in  $\mathbb{R}^n$ , then  $NO(A) \leq mass A$ .

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- If k > 0 is a positive integer, the multiply-by-k map f(T) = kT on the space of integral flat chains is an embedding with closed image.
- If T is a mod-k current, then T ≡ T<sub>Z</sub> (mod k) for some integral current T<sub>Z</sub>. Consequently, mod-k currents are a quotient of the integral currents.

Theorem If A is a mod-2 d-cycle in  $\mathbb{R}^n$ , then NO(A)  $\leq \max A$ .

#### Theorem

If A is a mod-2 d-cycle in  $\mathbb{R}^n$ , then  $NO(A) \lesssim mass A$ .

Strategy:

Find a mod-2 (d + 1)-chain such that  $A = \partial F$ .

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- ► Typically, F is non-orientable. Cut F into orientable pieces to get a lift F<sub>Z</sub> of F with integer coefficients.
- Then  $P = \partial F_{\mathbb{Z}}$  is a pseudo-orientation of A.
- ► The difference mass P mass A measures how much of F we had to cut.

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F is orientable, so A is orientable and NO(A) = mass(A).

# Example: the immersed Klein bottle

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so it is orientable!

# Results in low codimension

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#### Corollary (Federer)

If T is an integral (n-2)-cycle in  $\mathbb{R}^n$ , then FV(2T) = 2 FV(T).

# Results in low codimension

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Every (n-1)-cycle in  $\mathbb{R}^n$  is orientable, i.e., NO(A) = mass(A).

#### Corollary (Federer)

If T is an integral (n-2)-cycle in  $\mathbb{R}^n$ , then FV(2T) = 2 FV(T). What about higher codimensions?



Let A be a mod-2 cellular d-cycle of mass V



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- Orient the cubes at random to get  $F_{\mathbb{Z}}$
- $\partial F_{\mathbb{Z}}$  is a pseudo-orientation
- $\mathsf{NO}(A) \lesssim \mathsf{mass} \, \partial F_{\mathbb{Z}} \sim V^{(d+1)/d}$

# Bigger cubes



Total boundary:  $V^{(d+1)/d}$ 

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Total boundary:  $V^{(d+1)/d}$ 



Total boundary: much less











 $\sim$  V squares each with perimeter  $\sim$  1



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 $\sim$  V/2 squares each with perimeter  $\sim 2$ 

Sketch:

- ► Approximate A at ~ log V scales, then connect the approximations.
- We use cubes with total boundary  $\sim V$  at each scale.
- Since there are  $\sim \log V$  scales, we conclude:

#### Proposition (Guth-Y.)

If A is a cellular mod-2 cycle with volume V, then it has a pseudo-orientation P such that mass  $P \lesssim V \log V$ .

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- Choosing orientations randomly is wasteful when A is close to a plane
- But what if A is never close to a plane?

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How do we prove the proposition for sets that are close to fractals?

- Show that adding topological complexity adds extra area
- Prove the theorem when A has "low complexity"

#### Definition (David-Semmes)

A set  $E \subset \mathbb{R}^k$  is uniformly rectifiable if and only if E has a corona decomposition. (Roughly, for all but a few balls B, the intersection  $B \cap E$  is close to the graph of a Lipschitz function with small Lipschitz constant.)

#### Sketch of proof

#### Proposition

Every mod-2 cellular d-cycle A can be written as a sum

$$A = \sum_{i} A_{i}$$

of mod-2 cellular d-cycles with uniformly rectifiable support such that

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 mass  $A_i \leq C$  mass  $A$ .

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#### Proposition

Any mod-2 cellular d-cycle A with uniformly rectifiable support has a pseudo-orientation P with

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#### Open questions

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What does this tell us about the geometry of surfaces embedded in ℝ<sup>n</sup> by a bilipschitz map?